Chapter 7

SQL-99: Schema Definition, Basic Constraints, and Queries (from E&N and my editing)

Language

- Set operations to define and manipulate structure and constraints of database.
- All related to relational model
- Formal Language → Relational Algebra, Relational Calculus
- SQL → success in bussiness of DBMS tool

Data Definition, Constraints, and Schema Changes

- Used to CREATE, DROP, and ALTER the descriptions of the tables (relations) of a database

History

- SQL stand for Structured Query Language
- SQL is based on the Relational Tuple Calculus
- Evolved from SEQUEL: **S**tructured **E**nglish **QUE**ry **L**anguage - part of IBM's SYSTEM R, 1974
- SQL2 Supported by
  - ORACLE, SYBASE, INFORMIX,
  - IBM DB2, SQL SERVER, …
  - MS Access, MySQL, …
- SQL2 also called SQL/92 is evolved from SQL/86, SQL/89, all were ANSI & ISO standard
- Currently Working on SQL3/SQL-99 with OO Extensions
- Now – SQL is standard language for commercial relational DBMS

SQL Components

- Data Definition Language (DDL)
  - For External and Conceptual Schemas
  - Views - DDL for External Schemas
- Data Manipulation Language (DML)
  - Interactive DML Against External and Conceptual Schemas
  - Embedded DML in Host PLs (EQL, JDBC, etc.)
- Others
  - Integrity (Allowable Values/Referential)
  - Catalog and Dictionary Facilities
  - Transaction Control (Long-Duration and Batch)
  - Authorization (Who can Do What When)

DDL & DML

- Data Definition Language (DDL)
  - Defining the Relational Schema - Relations, Attributes, Domains - The Meta-Data

    CREATE TABLE Student:

    Name(CHAR(30)),SSN(CHAR(9)),GPA(FLOAT(2))

    CREATE TABLE Courses:

    Course#(CHAR(6)), Title(CHAR(20)), Descrip(CHAR(100)), Pcourse#(CHAR(6))
- Data Manipulation Language (DML)
  - Defining the Queries Against the Schema

    SELECT Name, SSN

    From Student

    Where GPA > 3.00

Look ...

- From Relational Model

STUDENT

Name	Ssn	GPA

COURSE

Course_num	Title	Description	Pre_Course_num

CREATE TABLE Student:
Name(CHAR(30)),SSN(CHAR(9)),GPA(FLOAT(2))
CREATE TABLE Courses:
Course#(CHAR(6)), Title(CHAR(20)), Descrip(CHAR(100)), Pcourse#(CHAR(6))

DDL: Data Definition Language

- A Pre-Defined set of Primitive Types
  - Numeric
  - Character-string
  - Bit-string
  - Additional Types
- Defining Domains
- Defining Schema
- Defining Tables
- Defining Views

DDL Primitive

- Numeric
  - INTEGER (or INT), SMALLINT
  - REAL, DOUBLE PRECISION
  - FLOAT(N) Floating Point with at Least N Digits
  - DECIMAL(P,D) (DEC(P,D) or NUMERIC(P,D)) have P Total Digits with D to Right of Decimal
- Note that INTs and REALs are Machine Dependent (Based on Hardware/OS Platform)

- Character-String
  - CHAR(N) or CHARACTER(N) – Fixed
  - VARCHAR(N), CHAR VARYING(N), or CHARACTER VARYING(N)
    Variable with at Most N Characters
- Bit-Strings
  - BIT(N) Fixed
- VARBIT(N) or BIT VARYING(N)
  - Variable with at Most N Bits

Additional ...

- Has DATE, TIME, and TIMESTAMP data types
  - DATE:

    Made up of year-month-day in the format yyyy-mm-dd
  - TIME:

    Made up of hour:minute:second in the format hh:mm:ss
  - TIME(i):

    Made up of hour:minute:second plus i additional digits specifying fractions of a second

    format is hh:mm:ss:ii...i
  - TIMESTAMP:

    Has both DATE and TIME components

Additional ...

- INTERVAL:
  - Specifies a relative value rather than an absolute value
  - Can be DAY/TIME intervals or YEAR/MONTH intervals
  - Can be positive or negative when added to or subtracted from an absolute value, the result is an absolute value

DDL Domains

- Domains are Similar in Concepts to Programming Language Type Definitions
- A Domain can be Defined as Follows:
  - **CREATE DOMAIN** CITY CHAR(15) **DEFAULT** '<Storrs>';
  - **CREATE DOMAIN** SSNFORMAT CHAR(9);
- Advantage of Using Domains
  - Changing a Domain Definition in One Place Changes it Consistently Everywhere it is Used
  - Default Values Can Be Defined for Domains
  - Constraints Can Be Defined for Domains

- A Domain is Dropped As Follows:
  - **DROP DOMAIN** CITY **RESTRICT;**
  - **DROP DOMAIN** SSNFORMAT **CASCADE;**
- Restrict:
  - Drop Operation Fails If the Domain is Used in Column Definitions
- Cascade:
  - Drop Operation Causes Columns to be Defined Directly on the Underlying Data Type

SQL - Relational Model

- Term Used

SQL	Formal Relational Model
Table	Relation
Row	Tuple
Column	Attribute

SQL Schema

- SQL Schema is identified by schema name and include authorization identifier.
- Schema elements: tables, attributes names, constraints, views, domains and other construct (such as authorization grant) that describe the schema
- System Administrator or DBA had privilege to create schemas
- Features that added to SQL2 & SQL-99

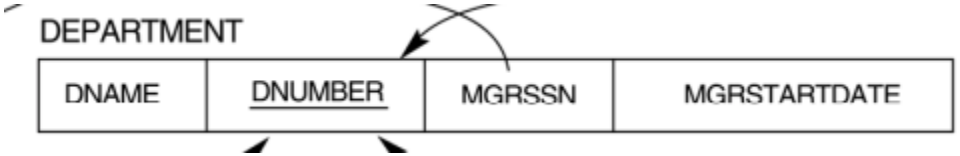
Creating/ Drop Schema

- Creating a Schema:
  CREATE SCHEMA MY_COMPANY AUTHORIZATION Dww;
  - Schema MY_COMPANY bas Been Created and is Owner by the User "Dww"
  - Tables can now be Created and Added to Schema
- Dropping a Schema:
  DROP SCHEMA MY_COMPANY RESTRICT;
  DROP SCHEMA MY_COMPANY CASCADE;
- Restrict:
  - Drop Operation Fails If Schema is Not Empty
- Cascade:
  - Drop Operation Removes Everything in the Schema

CREATE TABLE

- Specifies a new base relation by giving it a name, and specifying each of its attributes and their data types (INTEGER, FLOAT, DECIMAL(i,j), CHAR(n), VARCHAR(n))
- A constraint NOT NULL may be specified on an attribute

```
CREATE TABLE   DEPARTMENT
    (       DNAME          VARCHAR(10) NOT NULL,
            DNUMBER        INTEGER       NOT NULL,
            MGRSSN         CHAR(9),
            MGRSTARTDATE       CHAR(9)  );
```



CREATE TABLE

- In SQL2, can use the CREATE TABLE command for specifying the primary key attributes, secondary keys, and referential integrity constraints (foreign keys).
- Key attributes can be specified via the PRIMARY KEY and UNIQUE phrases

```
CREATE TABLE   DEPT
(    DNAME          VARCHAR(10) NOT NULL,
     DNUMBER        INTEGER       NOT NULL,
     MGRSSN         CHAR(9),
     MGRSTARTDATE       CHAR(9),
     PRIMARY KEY (DNUMBER),
     UNIQUE (DNAME),
     FOREIGN KEY (MGRSSN) REFERENCES EMP  );
```

DROP TABLE

- Used to remove a relation (base table) *and its definition*
- The relation can no longer be used in queries, updates, or any other commands since its description no longer exists
- Example:

```
DROP TABLE  DEPENDENT;
```

ALTER TABLE

- Used to add an attribute to one of the base relations
- The new attribute will have NULLs in all the tuples of the relation right after the command is executed; hence, the NOT NULL constraint is *not allowed* for such an attribute
- Example:

```
ALTER TABLE EMPLOYEE ADD JOB
VARCHAR(12);
```

- The database users must still enter a value for the new attribute JOB for each EMPLOYEE tuple. This can be done using the UPDATE command.

REFERENTIAL INTEGRITY OPTIONS

- We can specify RESTRICT, CASCADE, SET NULL or SET DEFAULT on referential integrity constraints (foreign keys)

```
CREATE TABLE   DEPT
 (   DNAME      VARCHAR(10)      NOT NULL,
     DNUMBER  INTEGER     NOT NULL,
     MGRSSN     CHAR(9),
     MGRSTARTDATE  CHAR(9),
     PRIMARY KEY (DNUMBER),
     UNIQUE (DNAME),
     FOREIGN KEY (MGRSSN) REFERENCES EMP
ON DELETE SET DEFAULT ON UPDATE CASCADE  );
```

REFERENTIAL INTEGRITY OPTIONS (continued)

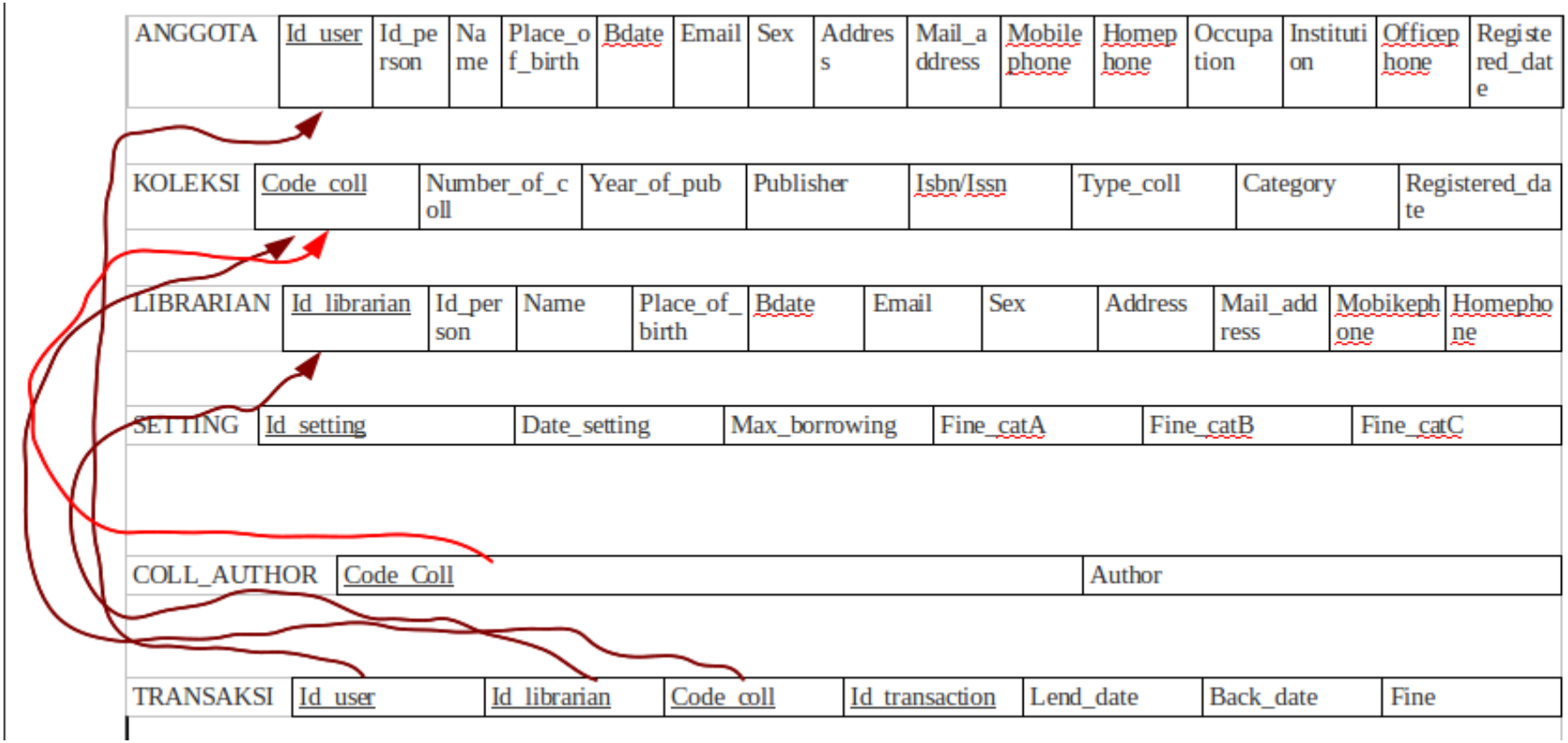
```
CREATE TABLE   EMP
    (       ENAME          VARCHAR(30) NOT NULL,
            ESSN   CHAR(9),
            BDATE DATE,
            DNO    INTEGER  DEFAULT 1,
            SUPERSSN      CHAR(9),
            PRIMARY KEY (ESSN),
            FOREIGN KEY (DNO) REFERENCES DEPT
          ON DELETE SET DEFAULT ON UPDATE CASCADE,
            FOREIGN KEY (SUPERSSN) REFERENCES EMP
          ON DELETE SET NULL ON UPDATE CASCADE  );
```

Implications of Drop/ Alter Table

- Possible Issues When you Drop or Alter a Table?
  - Views are Impacted - Portions (All?) of External Schema w.r.t. User Applications May No Longer be Available
  - User Applications May No Longer Execute
  - Applications that Utilize JDBC/ODBC to Access Conceptual Schema Directly May No Longer Work
  - Adding Columns via Alter Leads to …
    - Need to Update all Nulls with Actual Values
    - What if DB is Large?
    - Potential to Introduce Data Inconsistencies

Ex: Library Case :)

- Make DDL



- CREATE DATABASE LIBRARY;
- CREATE TABLE ANGGOTA (

  Id_user char(7),Id_person varchar(20),Name varchar(50),Place_ofB varchar(40), Bdate date,Email varchar(50),Sex char(1),Address varchar(100), Mail_address varchar(100), Mobilephone varchar(15), Homephone varchar(15), Occupation varchar(50), Institution varchar(50), Officephone varchar(15), Registered_date date,

  PRIMARY KEY Id_user);

Retrieval Queries in SQL

- SQL has one basic statement for retrieving information from a database; the SELECT statement
- Important distinction between SQL and the formal relational model; SQL allows a table (relation) to have two or more tuples that are identical in all their attribute values
- Hence, an SQL relation (table) is a *multi-set* (sometimes called a bag) of tuples; it *is not* a set of tuples
- SQL relations can be constrained to be sets by specifying PRIMARY KEY or UNIQUE attributes, or by using the DISTINCT option in a query

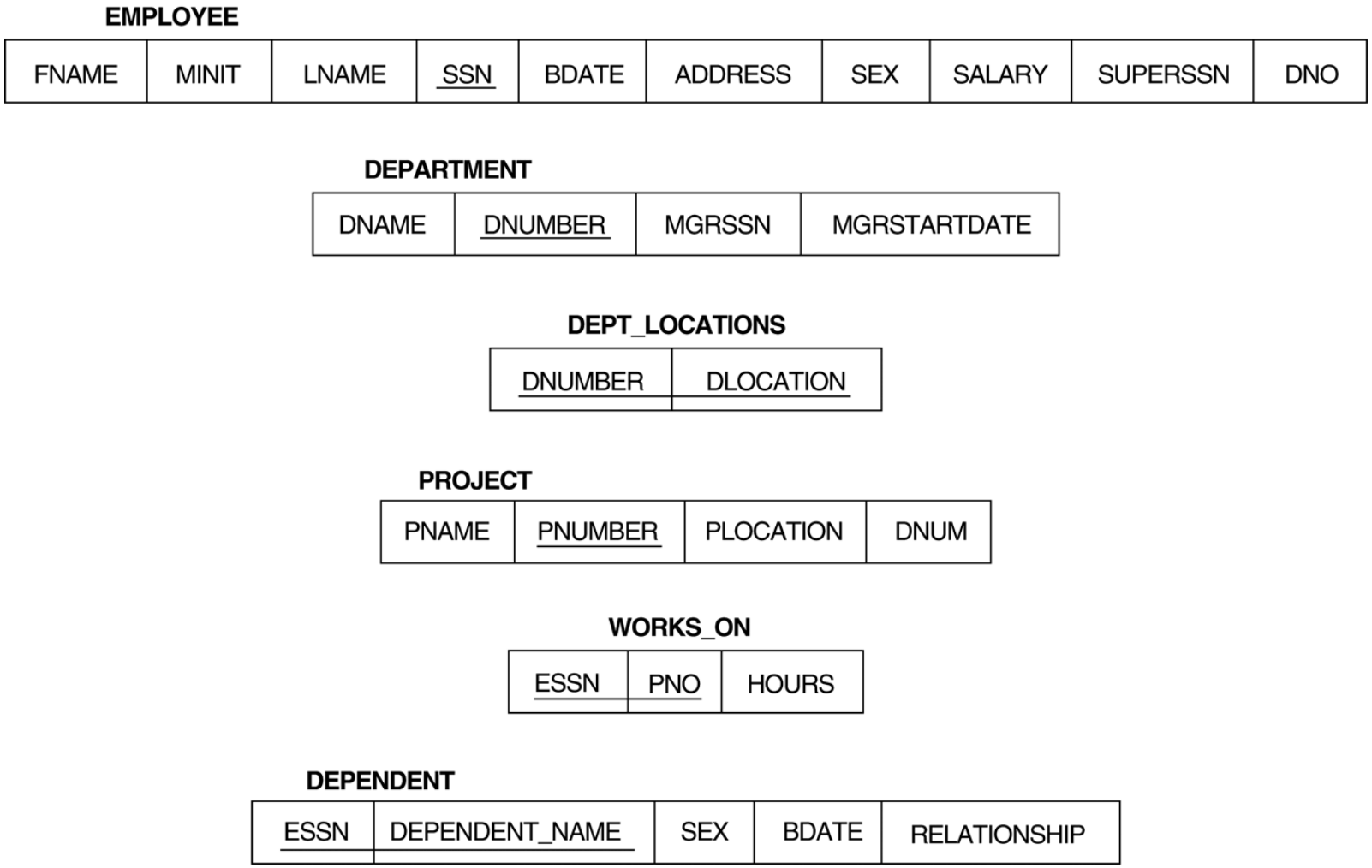
Retrieval Queries in SQL (cont.)

- Basic form of the SQL SELECT statement is called a *mapping* or a *SELECT-FROM-WHERE block*

```
SELECT   <attribute list>
FROM     <table list>
WHERE    <condition>
```

  - <attribute list> is a list of attribute names whose values are to be retrieved by the query
  - <table list> is a list of the relation names required to process the query
  - <condition> is a conditional (Boolean) expression that identifies the tuples to be retrieved by the query

Relational Database Schema



Populated Database

EMPLOYEE	FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO
	John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
	Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
	Alicia	J	Zelaya	999887777	1968-07-19	3321 Castle, Spring, TX	F	25000	987654321	4
	Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
	Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
	Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
	Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
	James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	null	1

DEPT_LOCATIONS	DNUMBER	DLOCATION
	1	Houston
	4	Stafford
	5	Bellaire
	5	Sugarland
	5	Houston

DEPARTMENT	DNAME	DNUMBER	MGRSSN	MGRSTARTDATE
	Research	5	333445555	1988-05-22
	Administration	4	987654321	1995-01-01
	Headquarters	1	888665555	1981-06-19

WORKS_ON	ESSN	PNO	HOURS
	123456789	1	32.5
	123456789	2	7.5
	666884444	3	40.0
	453453453	1	20.0
	453453453	2	20.0
	333445555	2	10.0
	333445555	3	10.0
	333445555	10	10.0
	333445555	20	10.0
	999887777	30	30.0
	999887777	10	10.0
	987987987	10	35.0
	987987987	30	5.0
	987654321	30	20.0
	987654321	20	15.0
	888665555	20	null

PROJECT	PNAME	PNUMBER	PLOCATION	DNUM
	ProductX	1	Bellaire	5
	ProductY	2	Sugarland	5
	ProductZ	3	Houston	5
	Computerization	10	Stafford	4
	Reorganization	20	Houston	1
	Newbenefits	30	Stafford	4

DEPENDENT	ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP
	333445555	Alice	F	1986-04-05	DAUGHTER
	333445555	Theodore	M	1983-10-25	SON
	333445555	Joy	F	1958-05-03	SPOUSE
	987654321	Abner	M	1942-02-28	SPOUSE
	123456789	Michael	M	1988-01-04	SON
	123456789	Alice	F	1988-12-30	DAUGHTER
	123456789	Elizabeth	F	1967-05-05	SPOUSE

Simple SQL Queries

- Basic SQL queries correspond to using the SELECT, PROJECT, and JOIN operations of the relational algebra
- All subsequent examples use the COMPANY database
- Example of a simple query on *one* relation
- Query 0: Retrieve the birthdate and address of the employee whose name is 'John B. Smith'.

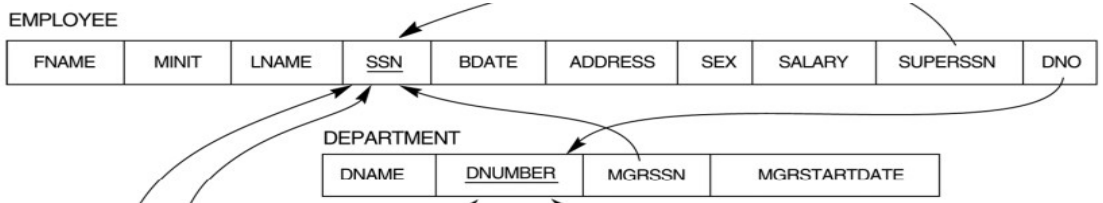
```
Q0: SELECT      BDATE, ADDRESS
    FROM        EMPLOYEE
    WHERE       FNAME='John' AND MINIT='B'
      AND       LNAME='Smith'
```

- Similar to a SELECT-PROJECT pair of relational algebra operations; the SELECT-clause specifies the *projection attributes* and the WHERE-clause specifies the *selection condition*
- However, the result of the query *may contain* duplicate tuples

Simple SQL Queries (cont.)

- Query 1: Retrieve the name and address of all employees who work for the 'Research' department.

```
Q1:     SELECT    FNAME, MINIT,LNAME, ADDRESS
        FROM      EMPLOYEE, DEPARTMENT
        WHERE     DNAME='Research' AND DNUMBER=DNO
```



  - Similar to a SELECT-PROJECT-JOIN sequence of relational algebra operations
  - (DNAME='Research') is a *selection condition* (corresponds to a SELECT operation in relational algebra)
  - (DNUMBER=DNO) is a *join condition* (corresponds to a JOIN operation in relational algebra)

Simple SQL Queries (cont.)

- Query 2: For every project located in 'Stafford', list the project number, the controlling department number, and the department manager's last name, address, and birthdate.

```
Q2:     SELECT      PNUMBER, DNUM, LNAME, BDATE,
        ADDRESS
          FROM      PROJECT, DEPARTMENT, EMPLOYEE
          WHERE     DNUM=DNUMBER AND MGRSSN=SSN
          AND       PLOCATION='Stafford'
```

  - In Q2, there are *two* join conditions
  - The join condition DNUM=DNUMBER relates a project to its controlling department
  - The join condition MGRSSN=SSN relates the controlling department to the employee who manages that department

Aliases, * and DISTINCT, Empty WHERE-clause

- In SQL, we can use the same name for two (or more) attributes as long as the attributes are in *different relations*

  A query that refers to two or more attributes with the same name must *qualify* the attribute name with the relation name by *prefixing* the relation name to the attribute name

Example:

- EMPLOYEE.LNAME, DEPARTMENT.DNAME

ALIASES

- Some queries need to refer to the same relation twice
- In this case, *aliases* are given to the relation name
- Query 8: For each employee, retrieve the employee's name, and the name of his or her immediate supervisor.

```
Q8: SELECT    E.FNAME, E.LNAME, S.FNAME,
                S.LNAME
      FROM    EMPLOYEE E S
      WHERE   E.SUPERSSN=S.SSN
```

  - In Q8, the alternate relation names E and S are called *aliases* or *tuple variables* for the EMPLOYEE relation
  - We can think of E and S as two *different copies* of EMPLOYEE; E represents employees in role of *supervisees* and S represents employees in role of *supervisors*

ALIASES (cont.)

- Aliasing can also be used in any SQL query for convenience
  Can also use the AS keyword to specify aliases

```
Q8:   SELECT   E.FNAME, E.LNAME, S.FNAME,
                S.LNAME
      FROM     EMPLOYEE AS E, EMPLOYEE AS S
      WHERE    E.SUPERSSN=S.SSN
```

UNSPECIFIED WHERE-clause

- A *missing WHERE-clause* indicates no condition; hence, *all tuples* of the relations in the FROM-clause are selected
- This is equivalent to the condition WHERE TRUE
- Query 9: Retrieve the SSN values for all employees.

```
Q9:     SELECT    SSN
  FROM  EMPLOYEE
```

- If more than one relation is specified in the FROM-clause *and* there is no join condition, then the *CARTESIAN PRODUCT* of tuples is selected

UNSPECIFIED WHERE-clause (cont.)

- Example:

```
Q10:    SELECT    SSN, DNAME
        FROM      EMPLOYEE, DEPARTMENT
```

  - It is extremely important not to overlook specifying any selection and join conditions in the WHERE-clause; otherwise, incorrect and very large relations may result

USE OF *

- To retrieve all the attribute values of the selected tuples, a * is used, which stands for *all the attributes*
  Examples:

```
Q1C:    SELECT    *
        FROM      EMPLOYEE
        WHERE     DNO=5

Q1D:    SELECT    *
        FROM      EMPLOYEE, DEPARTMENT
        WHERE     DNAME='Research' AND
                        DNO=DNUMBER
```

USE OF DISTINCT

- SQL does not treat a relation as a set; *duplicate tuples can appear*
- To eliminate duplicate tuples in a query result, the keyword **DISTINCT** is used
- For example, the result of Q11 may have duplicate SALARY values whereas Q11A does not have any duplicate values

```
Q11:    SELECT    SALARY
        FROM      EMPLOYEE
Q11A:   SELECT    DISTINCT SALARY
        FROM      EMPLOYEE
```

SET OPERATIONS

- SQL has directly incorporated some set operations
- There is a union operation (**UNION**), and in *some versions* of SQL there are set difference (**MINUS**) and intersection (**INTERSECT**) operations
- The resulting relations of these set operations are sets of tuples; *duplicate tuples are eliminated from the result*
- The set operations apply only to *union compatible relations* ; the two relations must have the same attributes and the attributes must appear in the same order

SET OPERATIONS (cont.)_

- Query 4: Make a list of all project numbers for projects that involve an employee whose last name is 'Smith' as a worker or as a manager of the department that controls the project.

```
Q4: (SELECT  PNAME
     FROM        PROJECT, DEPARTMENT, EMPLOYEE
     WHERE       DNUM=DNUMBER AND
MGRSSN=SSN       AND          LNAME='Smith')
     UNION       (SELECT  PNAME
     FROM        PROJECT, WORKS_ON, EMPLOYEE
     WHERE       PNUMBER=PNO AND ESSN=SSN AND
                          LNAME='Smith')
```

NESTING OF QUERIES

- A complete SELECT query, called a *nested query* , can be specified within the WHERE-clause of another query, called the *outer query*
- Many of the previous queries can be specified in an alternative form using nesting
- Query 1: Retrieve the name and address of all employees who work for the 'Research' department.

```
Q1: SELECT      FNAME, LNAME, ADDRESS
    FROM        EMPLOYEE
    WHERE       DNO IN  (SELECT  DNUMBER
                FROM        DEPARTMENT
                WHERE       DNAME='Research' )
```

NESTING OF QUERIES (cont.)

- The nested query *selects the number of the 'Research' department*
- The *outer query select an EMPLOYEE tuple* if its DNO value is in the result of either nested query
- The comparison operator **IN** compares a value v with a set (or multi-set) of values V, and evaluates to **TRUE** if v is one of the elements in V
- In general, we can have several levels of nested queries
- A reference to an *unqualified attribute* refers to the relation declared in the *innermost nested query*
- In this example, the nested query is *not correlated* with the outer query

CORRELATED NESTED QUERIES

- If a condition in the WHERE-clause of a *nested query* references an attribute of a relation declared in the *outer query* , the two queries are said to be *correlated*
- The result of a correlated nested query is *different for each tuple (or combination of tuples) of the relation(s) the outer query*
- Query 12: Retrieve the name of each employee who has a dependent with the same first name as the employee.

```
Q12: SELECT     E.FNAME, E.LNAME
     FROM       EMPLOYEE AS E
     WHERE      E.SSN IN (SELECT    ESSN
                FROM  DEPENDENT
                WHERE       ESSN=E.SSN AND

E.FNAME=DEPENDENT_NAME)
```

CORRELATED NESTED QUERIES (cont.)

- In Q12, the nested query has a different result *for each tuple* in the outer query
- A query written with nested SELECT... FROM... WHERE... blocks and using the = or IN comparison operators can ***always*** be expressed as a single block query. For example, Q12 may be written as in Q12A

```
Q12A:   SELECT    E.FNAME, E.LNAME
        FROM      EMPLOYEE E, DEPENDENT D
        WHERE     E.SSN=D.ESSN AND
                  E.FNAME=D.DEPENDENT_NAME
```

- The original SQL as specified for SYSTEM R also had a **CONTAINS** comparison operator, which is used in conjunction with nested correlated queries
- This operator was dropped from the language, possibly because of the difficulty in implementing it efficiently

CORRELATED NESTED QUERIES (cont.)

- Most implementations of SQL *do not* have this operator
- The CONTAINS operator compares two *sets of values* , and returns TRUE if one set contains all values in the other set (reminiscent of the *division* operation of algebra).
  - Query 3: Retrieve the name of each employee who works on *all* the projects controlled by department number 5.

```
Q3:   SELECT  FNAME, LNAME
      FROM    EMPLOYEE
      WHERE  ( (SELECT          PNO
                FROM  WORKS_ON
                WHERE           SSN=ESSN)
                CONTAINS
               (SELECT          PNUMBER
                FROM  PROJECT
                WHERE           DNUM=5) )
```

CORRELATED NESTED QUERIES (cont.)

- In Q3, the second nested query, which is not correlated with the outer query, retrieves the project numbers of all projects controlled by department 5
- The first nested query, which is correlated, retrieves the project numbers on which the employee works, which is different *for each employee tuple* because of the correlation

THE EXISTS FUNCTION

- EXISTS is used to check whether the result of a correlated nested query is empty (contains no tuples) or not
- We can formulate Query 12 in an alternative form that uses EXISTS as Q12B below

THE EXISTS FUNCTION (cont.)

- Query 12: Retrieve the name of each employee who has a dependent with the same first name as the employee.

```
Q12B:  SELECT   FNAME, LNAME
       FROM     EMPLOYEE
       WHERE    EXISTS   (SELECT   *
                FROM     DEPENDENT
                WHERE    SSN=ESSN AND

FNAME=DEPENDENT_NAME)
```

THE EXISTS FUNCTION (cont.)

- Query 6: Retrieve the names of employees who have no dependents.

```
Q6:     SELECT   FNAME, LNAME
        FROM     EMPLOYEE
        WHERE    NOT EXISTS  (SELECT  *
                 FROM  DEPENDENT
                 WHERE SSN=ESSN)
```

  - In Q6, the correlated nested query retrieves all DEPENDENT tuples related to an EMPLOYEE tuple. If *none exist* , the EMPLOYEE tuple is selected
  - EXISTS is necessary for the expressive power of SQL

EXPLICIT SETS

- It is also possible to use an explicit (enumerated) set of values in the WHERE-clause rather than a nested query
- Query 13: Retrieve the social security numbers of all employees who work on project number 1, 2, or 3.

```
Q13:    SELECT    DISTINCT ESSN
        FROM      WORKS_ON
        WHERE     PNO IN  (1, 2, 3)
```

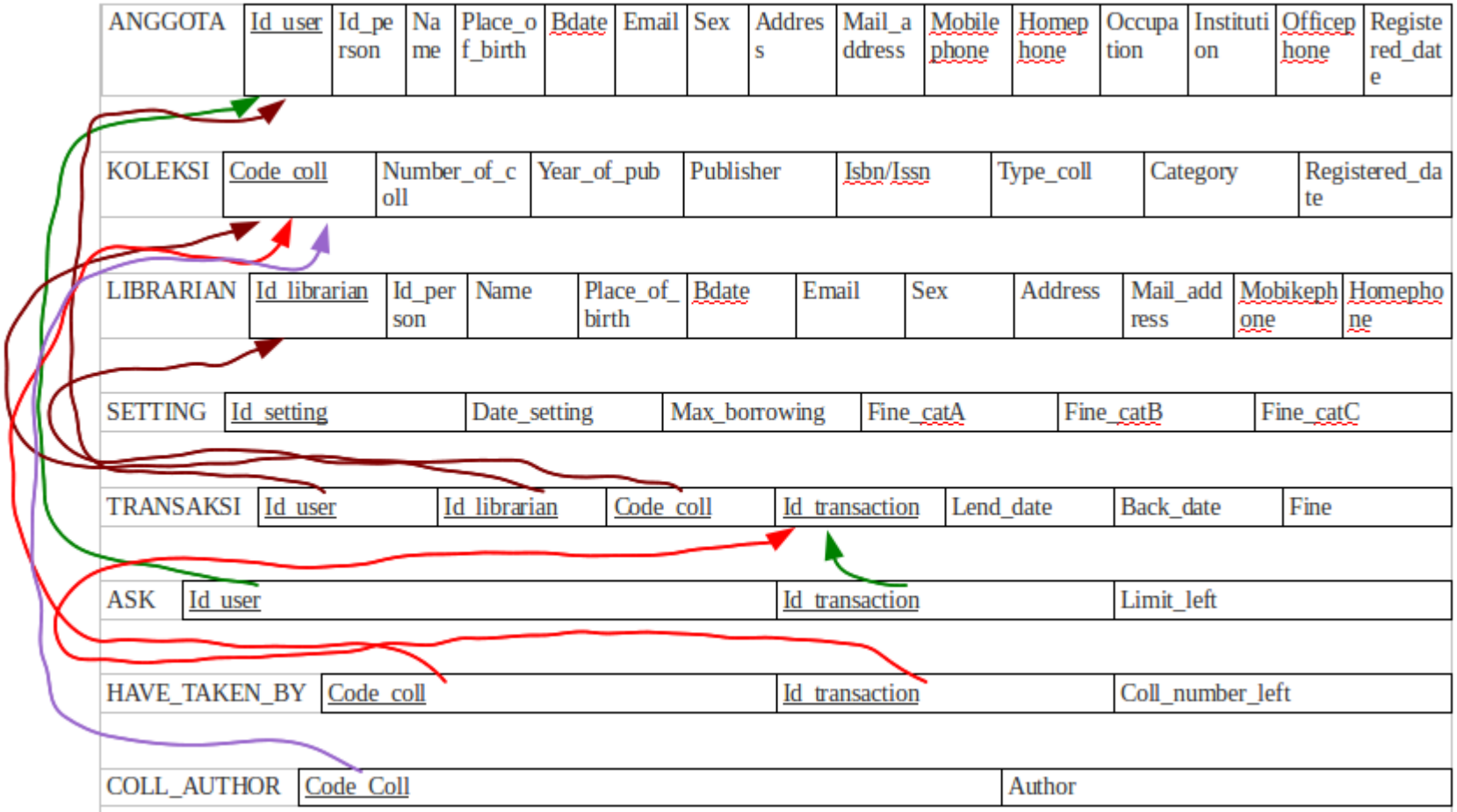
NULLS IN SQL QUERIES

- SQL allows queries that check if a value is NULL (missing or undefined or not applicable)
- SQL uses IS or IS NOT to compare NULLs because it considers each NULL value distinct from other NULL values, so equality comparison is not appropriate .
- Query 14: Retrieve the names of all employees who do not have supervisors.

```
Q14:    SELECT    FNAME, LNAME
        FROM      EMPLOYEE
        WHERE     SUPERSSN  IS  NULL
```

Note: If a join condition is specified, tuples with NULL values for the join attributes are not included in the result

Ex: Library Case



- Q0: Retrieve semua ID anggota yang pernah meminjam

```
SELECT Id_user
FROM TRANSAKSI
```

- Q1: Retrieve semua nama anggota yang pernah meminjam

```
SELECT Name
FROM ANGGOTA AS A
WHERE A.Id_user IN (
    SELECT DISTICT Id_user
    FROM TRANSAKSI AS T
    WHERE T.Id_user = A.Id_user)
```

- Q2: Retrieve nama peminjam dan kode koleksi yang memiliki Issn/Isbn dan belum dikembalikan
- Q3: Retrieve nama librarian yang melayani transaksi dengan denda > 10.000
- Q4: Retrieve koleksi yang semuanya terpinjam
- Q5: Retrieve nama anggota dan koleksi yang dikembalikan yang tidak mendapat denda hari ini

- Retrieve nama peminjam dan kode koleksi yang memiliki Issn/Isbn dan belum dikembalikan
- SELECT A.Name, T.Code_coll

  FROM ANGGOTA A, TRANSAKSI T, KOLEKSI K

  WHERE T.Id_user = A.Id_user AND T.Code_coll = K.Code_coll AND K.Issn/Isbn is NOT NULL AND T.Back_date is NULL

- Retrieve nama librarian yang melayani transaksi dengan denda > 10.000
- SELECT L.Name

  FROM LIBRARIAN L, TRANSAKSI T

  WHERE L.Id_librarian = T.Id_Librarian AND T.Fine > 10.000

- Retrieve koleksi yang semuanya terpinjam
- SELECT K.*

  FROM KOLEKSI K, HAVE_TAKEN_BY H

  WHERE K.Code_coll = H.Code_coll AND H.Code_number_left = “ ”

- Retrieve nama anggota dan kode koleksi yang dikembalikan yang tidak mendapat denda hari ini
- SELECT A.Name, T.Code_coll

  FROM ANGGOTA A, TRANSAKSI T

  WHERE T.Back_date = “11/4/2010” AND T.Fine = “ ” AND T.Id_user = A.Id_user

Joined Relations Feature in SQL2

- Can specify a "joined relation" in the FROM-clause
- Looks like any other relation but is the result of a join
- Allows the user to specify different types of joins (regular "theta" JOIN, NATURAL JOIN, LEFT OUTER JOIN, RIGHT OUTER JOIN, CROSS JOIN, etc)

Joined Relations Feature in SQL2 (cont.)

- Examples:

```
Q8: SELECT    E.FNAME, E.LNAME, S.FNAME, S.LNAME
    FROM      EMPLOYEE E S
    WHERE     E.SUPERSSN=S.SSN
```

can be written as:

```
Q8: SELECT    E.FNAME, E.LNAME, S.FNAME, S.LNAME
    FROM      (EMPLOYEE E LEFT OUTER JOIN
              EMPLOYEES ON  E.SUPERSSN=S.SSN)

Q1: SELECT    FNAME, LNAME, ADDRESS
    FROM      EMPLOYEE, DEPARTMENT
    WHERE     DNAME='Research' AND
DNUMBER=DNO
```

Joined Relations Feature in SQL2 (cont.)

- could be written as:

```
Q1: SELECT      FNAME, LNAME, ADDRESS
    FROM        (EMPLOYEE JOIN DEPARTMENT
                ON DNUMBER=DNO)
    WHERE       DNAME='Research'
```

or as:

```
Q1: SELECT      FNAME, LNAME, ADDRESS
    FROM        (EMPLOYEE NATURAL JOIN
                    DEPARTMENT AS
DEPT(DNAME, DNO,                    MSSN,
MSDATE)
    WHERE       DNAME='Research'
```

Joined Relations Feature in SQL2 (cont.)

- Another Example;
  - Q2 could be written as follows; this illustrates multiple joins in the joined tables

```
Q2:   SELECT    PNUMBER, DNUM,
      LNAME, BDATE, ADDRESS
      FROM      (PROJECT JOIN
                DEPARTMENT ON
                DNUM=DNUMBER) JOIN
                EMPLOYEE ON
                MGRSSN=SSN) )
      WHERE      PLOCATION='Stafford'
```

AGGREGATE FUNCTIONS

- Include **COUNT**, **SUM**, **MAX**, **MIN**, and **AVG**
- Query 15: Find the maximum salary, the minimum salary, and the average salary among all employees.

```
Q15:    SELECT    MAX(SALARY),
                  MIN(SALARY), AVG(SALARY)
        FROM      EMPLOYEE
```

  - Some SQL implementations *may not allow more than one function* in the SELECT-clause

AGGREGATE FUNCTIONS (cont.)

- Query 16: Find the maximum salary, the minimum salary, and the average salary among employees who work for the 'Research' department.

```
Q16: SELECT   MAX(SALARY), MIN(SALARY),
                    AVG(SALARY)
     FROM     EMPLOYEE, DEPARTMENT
     WHERE    DNO=DNUMBER AND
              DNAME='Research'
```

AGGREGATE FUNCTIONS (cont.)

- Queries 17 and 18: Retrieve the total number of employees in the company (Q17), and the number of employees in the 'Research' department (Q18).

```
Q17:    SELECT    COUNT (*)
        FROM      EMPLOYEE

Q18:    SELECT    COUNT (*)
        FROM      EMPLOYEE E S,   DEPARTMENT
        WHERE     DNO=DNUMBER AND
DNAME='Research' AND   E.SUPERSSN = E.SSN
AND S.Name = DEWi
```

GROUPING

- In many cases, we want to apply the aggregate functions *to subgroups of tuples in a relation*
- Each subgroup of tuples consists of the set of tuples that have *the same value* for the *grouping attribute(s)*
- The function is applied to each subgroup independently
- SQL has a **GROUP BY**-clause for specifying the grouping attributes, which *must also appear in the SELECT-clause*

GROUPING (cont.)

- Query 20: For each department, retrieve the department number, the number of employees in the department, and their average salary.

```
Q20:    SELECT    DNO, COUNT (*), AVG
(SALARY)
        FROM  EMPLOYEE
        GROUP BY    DNO
```

  - In Q20, the EMPLOYEE tuples are divided into groups--each group having the same value for the grouping attribute DNO
  - The COUNT and AVG functions are applied to each such group of tuples separately
  - The SELECT-clause includes only the grouping attribute and the functions to be applied on each group of tuples
  - A join condition can be used in conjunction with grouping

GROUPING (cont.)

- Query 21: For each project, retrieve the project number, project name, and the number of employees who work on that project.

```
Q21:    SELECT    PNUMBER, PNAME, COUNT (*)
        FROM      PROJECT, WORKS_ON
        WHERE     PNUMBER=PNO
        GROUP BY  PNUMBER, PNAME
```

  - In this case, the grouping and functions are applied *after* the joining of the two relations

THE HAVING-CLAUSE

- Sometimes we want to retrieve the values of these functions for only those *groups that satisfy certain conditions*
- The HAVING-clause is used for specifying a selection condition on groups (rather than on individual tuples)

THE HAVING-CLAUSE (cont.)

- Query 22: For each project *on which more than two employees work* , retrieve the project number, project name, and the number of employees who work on that project.

```
Q22:    SELECT    PNUMBER, PNAME, COUNT (*)
        FROM      PROJECT, WORKS_ON
        WHERE     PNUMBER=PNO
        GROUP BY  PNUMBER, PNAME
        HAVING    COUNT (*) > 2
```

SUBSTRING COMPARISON

- The LIKE comparison operator is used to compare partial strings
- Two reserved characters are used: '%' (or '*' in some implementations) replaces an arbitrary number of characters, and '_' replaces a single arbitrary character

SUBSTRING COMPARISON (cont.)

- Query 25: Retrieve all employees whose address is in Houston, Texas. Here, the value of the ADDRESS attribute must contain the substring 'Houston,TX'.

```
Q25:    SELECT    FNAME, LNAME
        FROM          EMPLOYEE
        WHERE     ADDRESS LIKE
                  '%Houston,TX%'
```

SUBSTRING COMPARISON (cont.)

- Query 26: Retrieve all employees who were born during the 1950s. Here, '5' must be the 8th character of the string (according to our format for date), so the BDATE value is '_____5_', with each underscore as a place holder for a single arbitrary character.

```
Q26:    SELECT    FNAME, LNAME
        FROM      EMPLOYEE
        WHERE     BDATE LIKE      '_______5_'
```

- The LIKE operator allows us to get around the fact that each value is considered atomic and indivisible; hence, in SQL, character string attribute values are not atomic

ARITHMETIC OPERATIONS

- The standard arithmetic operators '+', '-'. '*', and '/' (for addition, subtraction, multiplication, and division, respectively) can be applied to numeric values in an SQL query result
- Query 27: Show the effect of giving all employees who work on the 'ProductX' project a 10% raise.

```
Q27:    SELECT    FNAME, LNAME, 1.1*SALARY
                  FROM  EMPLOYEE, WORKS_ON,
PROJECT
        WHERE     SSN=ESSN AND PNO=PNUMBER AND
                          PNAME='ProductX'
```

ORDER BY

- The **ORDER BY** clause is used to sort the tuples in a query result based on the values of some attribute(s)
- Query 28: Retrieve a list of employees and the projects each works in, ordered by the employee's department, and within each department ordered alphabetically by employee last name.

```
Q28:    SELECT    DNAME, LNAME, FNAME, PNAME
    FROM  DEPARTMENT, EMPLOYEE,  WORKS_ON, PROJECT
    WHERE    DNUMBER=DNO AND SSN=ESSN   AND
PNO=PNUMBER
    ORDER BY    DNAME, LNAME
```

ORDER BY (cont.)

- The default order is in ascending order of values
- We can specify the keyword **DESC** if we want a descending order; the keyword **ASC** can be used to explicitly specify ascending order, even though it is the default

Summary of SQL Queries

- A query in SQL can consist of up to six clauses, but only the first two, SELECT and FROM, are mandatory. The clauses are specified in the following order:

```
SELECT   <attribute list>
FROM     <table list>
[WHERE   <condition>]
[GROUP BY <grouping attribute(s)>]
[HAVING <group condition>]
[ORDER BY <attribute list>]
```

Summary of SQL Queries (cont.)

- The SELECT-clause lists the attributes or functions to be retrieved
- The FROM-clause specifies all relations (or aliases) needed in the query but not those needed in nested queries
- The WHERE-clause specifies the conditions for selection and join of tuples from the relations specified in the FROM-clause
- GROUP BY specifies grouping attributes
- HAVING specifies a condition for selection of groups
- ORDER BY specifies an order for displaying the result of a query
- A query is evaluated by first applying the WHERE-clause, then GROUP BY and HAVING, and finally the SELECT-clause

Specifying Updates in SQL

- There are three SQL commands to modify the database; INSERT, DELETE, and UPDATE

INSERT

- In its simplest form, it is used to add one or more tuples to a relation
- Attribute values should be listed in the same order as the attributes were specified in the CREATE TABLE command

INSERT (cont.)

- Example:

```
U1:  INSERT INTO  EMPLOYEE
     VALUES (`Richard',,'Marini', '653298653', '30-
DEC-52',
     '98 Oak Forest,Katy,TX', 'M', 37000,'987654321',4
 )
```

- An alternate form of INSERT specifies explicitly the attribute names that correspond to the values in the new tuple
- Attributes with NULL values can be left out
- Example: Insert a tuple for a new EMPLOYEE for whom we only know the FNAME, LNAME, and SSN attributes.

```
U1A:   INSERT INTO EMPLOYEE (FNAME, LNAME, SSN)
       VALUES ('Richard', 'Marini', '653298653')
```

INSERT (cont.)

- Important Note: Only the constraints specified in the DDL commands are automatically enforced by the DBMS when updates are applied to the database
- Another variation of INSERT allows insertion of *multiple tuples* resulting from a query into a relation

INSERT (cont.)

- Example: Suppose we want to create a temporary table that has the name, number of employees, and total salaries for each department. A table DEPTS_INFO is created by U3A, and is loaded with the summary information retrieved from the database by the query in U3B.

```
U3A:    CREATE TABLE  DEPTS_INFO
                (DEPT_NAME   VARCHAR(10),
                 NO_OF_EMPS  INTEGER,
                 TOTAL_SAL   INTEGER);

U3B:    INSERT INTO   DEPTS_INFO (DEPT_NAME,
                      NO_OF_EMPS, TOTAL_SAL)
        SELECT        DNAME, COUNT (*), SUM
(SALARY)
        FROM          DEPARTMENT, EMPLOYEE
        WHERE         DNUMBER=DNO
        GROUP BY      DNAME ;
```

INSERT (cont.)

- Note: The DEPTS_INFO table may not be up-to-date if we change the tuples in either the DEPARTMENT or the EMPLOYEE relations *after* issuing U3B. We have to create a view (see later) to keep such a table up to date.

DELETE

- Removes tuples from a relation
- Includes a WHERE-clause to select the tuples to be deleted
- Tuples are deleted from only *one table* at a time (unless CASCADE is specified on a referential integrity constraint)
- A missing WHERE-clause specifies that *all tuples* in the relation are to be deleted; the table then becomes an empty table
- The number of tuples deleted depends on the number of tuples in the relation that satisfy the WHERE-clause
- Referential integrity should be enforced

DELETE (cont.)

- Examples:

```
U4A:    DELETE FROM    EMPLOYEE
        WHERE    LNAME='Brown'

U4B:    DELETE FROM    EMPLOYEE
        WHERE    SSN='123456789'

U4C:    DELETE FROM    EMPLOYEE
        WHERE    DNO  IN          (SELECT
DNUMBER
        FROM     DEPARTMENT
        WHERE    DNAME='Research')

U4D:    DELETE FROM    EMPLOYEE
```

UPDATE

- Used to modify attribute values of one or more selected tuples
- A WHERE-clause selects the tuples to be modified
- An additional SET-clause specifies the attributes to be modified and their new values
- Each command modifies tuples *in the same relation*
- Referential integrity should be enforced

UPDATE (cont.)

- Example: Change the location and controlling department number of project number 10 to 'Bellaire' and 5, respectively.

```
U5: UPDATE    PROJECT
    SET       PLOCATION = 'Bellaire', DNUM
= 5
    WHERE     PNUMBER=10
```

UPDATE (cont.)

- Example: Give all employees in the 'Research' department a 10% raise in salary.

```
U6:  UPDATE     EMPLOYEE
     SET        SALARY = SALARY *1.1
     WHERE      DNO  IN (SELECT     DNUMBER
                  FROM       DEPARTMENT
                  WHERE      DNAME='Research')
```

- In this request, the modified SALARY value depends on the original SALARY value in each tuple
- The reference to the SALARY attribute on the right of = refers to the old SALARY value before modification
- The reference to the SALARY attribute on the left of = refers to the new SALARY value after modification

Views in SQL

- A view is a “virtual” table that is derived from other tables
- Allows for limited update operations (since the table may not physically be stored)
- Allows full query operations
- A convenience for expressing certain operations

Specification of VIEWS

- SQL command: `CREATE VIEW`
  - a table (view) name
  - a possible list of attribute names (for example, when arithmetic operations are specified or when we want the names to be different from the attributes in the base relations)
  - a query to specify the table contents

VIEWS Ex..

- Specify a different WORKS_ON table

```
CREATE VIEW WORKS_ON_NEW AS
SELECT FNAME, LNAME, PNAME, HOURS
FROM EMPLOYEE, PROJECT, WORKS_ON
     WHERE SSN=ESSN AND PNO=PNUMBER
GROUP BY PNAME;
```

Using Virtual Table

- We can specify SQL queries on a newly create table (view):

```
SELECT FNAME, LNAME FROM
WORKS_ON_NEW
WHERE PNAME='Seena';
```

- When no longer needed, a view can be dropped:

```
DROP WORKS_ON_NEW;
```

Efficient VIEWS Implementation

- Query modification: present the view query in terms of a query on the underlying base tables
  - Disadvantage: inefficient for views defined via complex queries (especially if additional queries are to be applied to the view within a short time period)

- View materialization: involves physically creating and keeping a temporary table
  - assumption: other queries on the view will follow
  - concerns: maintaining correspondence between the base table and the view when the base table is updated
  - strategy: incremental update

VIEW Update

- Update on a single view without aggregate operations: update may map to an update on the underlying base table
- Views involving joins: an update *may* map to an update on the underlying base relations
  - not always possible

Un-Updatable VIEWS

- Views defined using groups and aggregate functions are not updateable
- Views defined on multiple tables using joins are generally not updateable
- `WITH CHECK OPTION`: must be added to the definition of a view if the view is to be updated
  - to allow check for updatability and to plan for an execution strategy