Chapter 5

Relational Database Design by ER- and EER-to-Relational Mapping (from E&N and my editing)

Outline

- ER-to-Relational Mapping Algorithm

  Step 1: Mapping of Regular Entity Types
  Step 2: Mapping of Weak Entity Types
  Step 3: Mapping of Binary 1:1 Relation Types
  Step 4: Mapping of Binary 1:N Relationship Types.
  Step 5: Mapping of Binary M:N Relationship Types.
  Step 6: Mapping of Multivalued attributes.
  Step 7: Mapping of N-ary Relationship Types.

- Mapping EER Model Constructs to Relations

  Step 8: Options for Mapping Specialization or Generalization.
  Step 9: Mapping of Union Types (Categories).

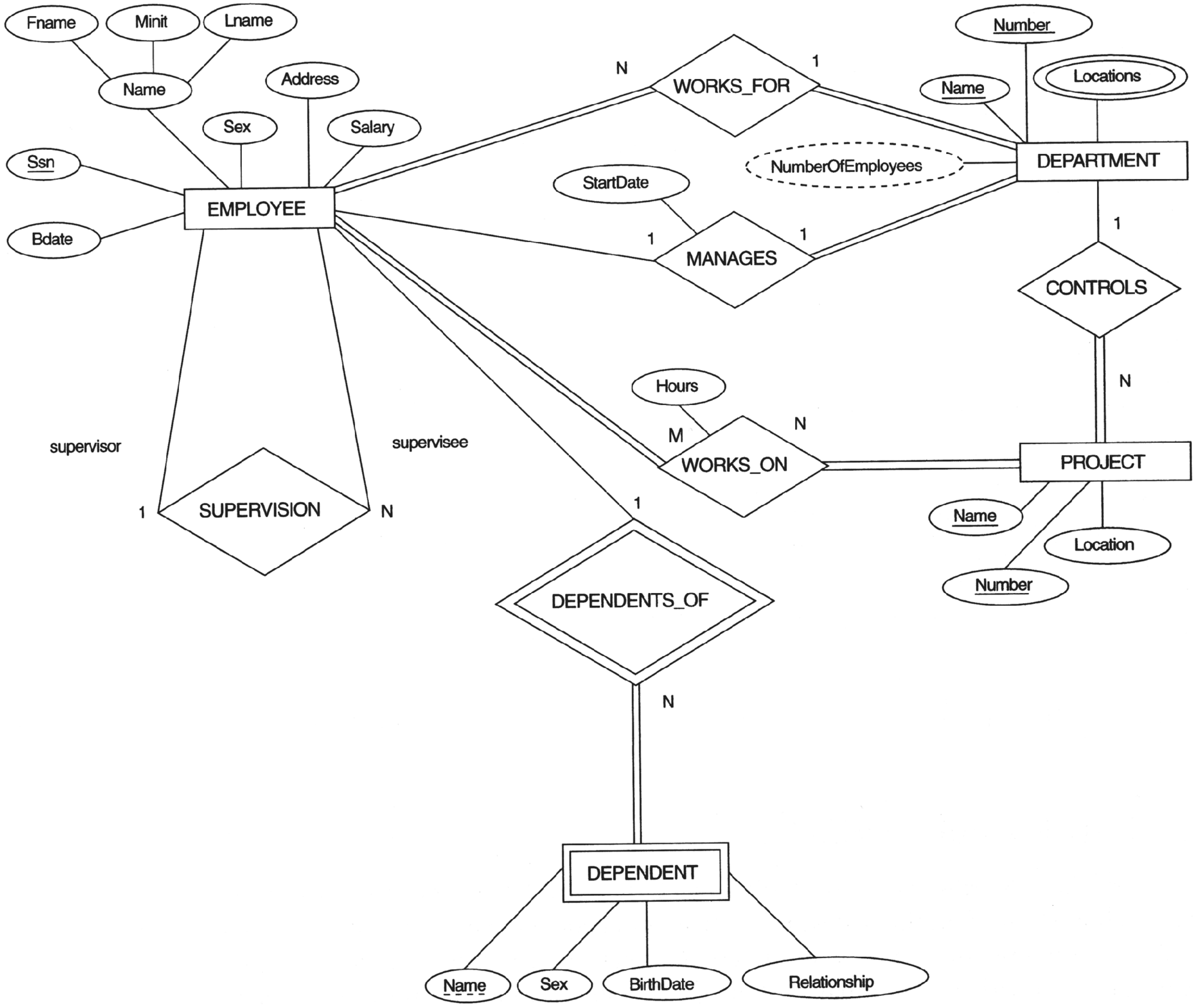
A bit explaination/req of COMPANY case

- Semua EMPLOYEE harus bekerja untuk hanya dalam satu DEPARTMENT
- DEPARTMENT dapat mengontrol beberapa PROJECT
- EMPLOYEE dapat bekerja pada lebih dari satu PROJECT
- DEPARTMENT dimanage oleh tepat satu MANAGER

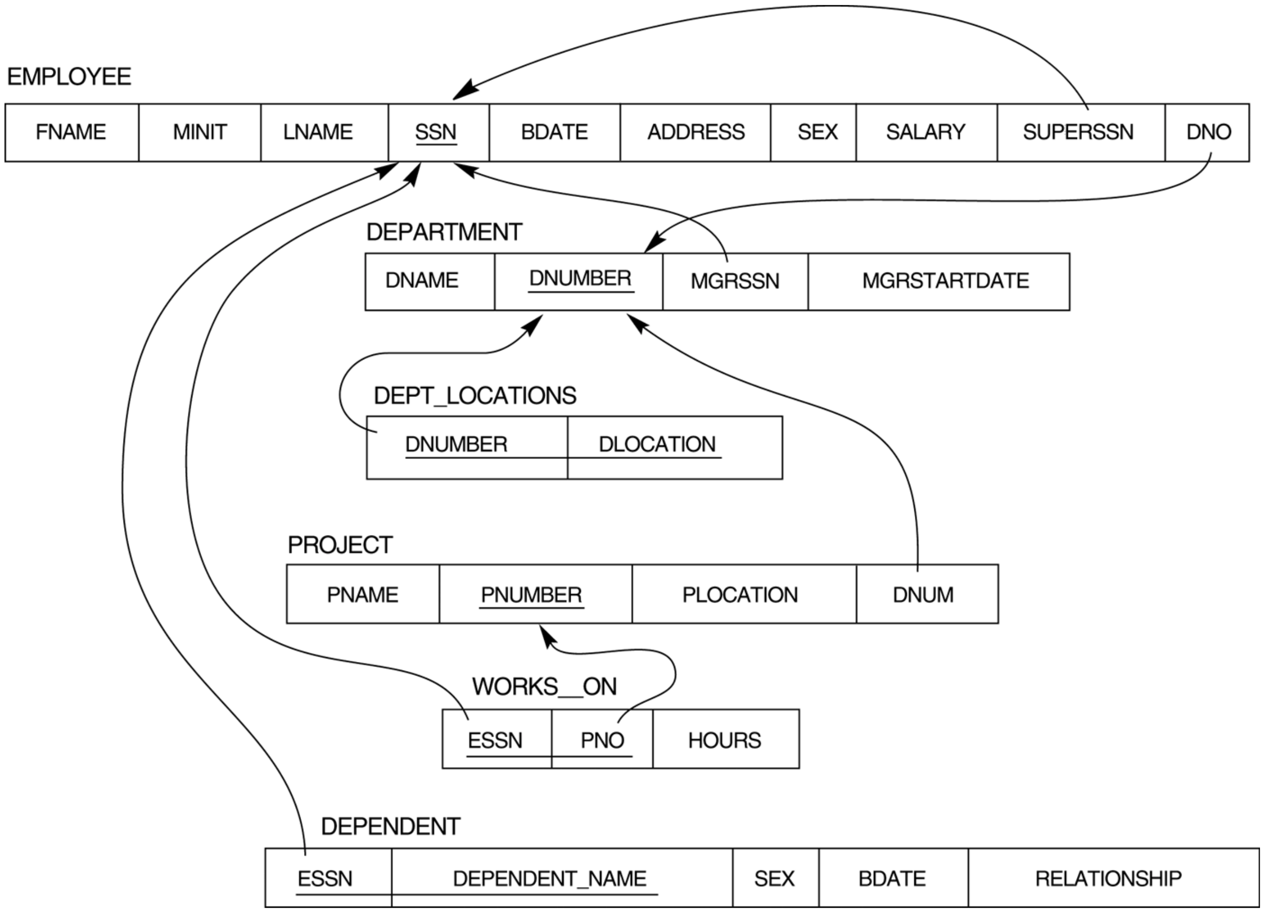
- MANAGER termasuk EMPLOYEE
- EMPLOYEE ada yang menjadi supervisor terhadap EMPLOYEE yang lain
- EMPLOYEE memiliki DEPENDENT (person yang dependent kepada EMPLOYEE)

- Note: In the real cases, realworld or miniworld, U as database designer must have ability to analys the requirements
- How?practice and improve Ur logic knowledge makes better

The ER conceptual schema diagram for the COMPANY database.



Result of mapping the COMPANY ER schema into a relational schema.



EMPLOYEE

FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO

DEPARTMENT

DNAME	DNUMBER	MGRSSN	MGRSTARTDATE

DEPT_LOCATIONS

DNUMBER	DLOCATION

PROJECT

PNAME	PNUMBER	PLOCATION	DNUM

WORKS__ON

ESSN	PNO	HOURS

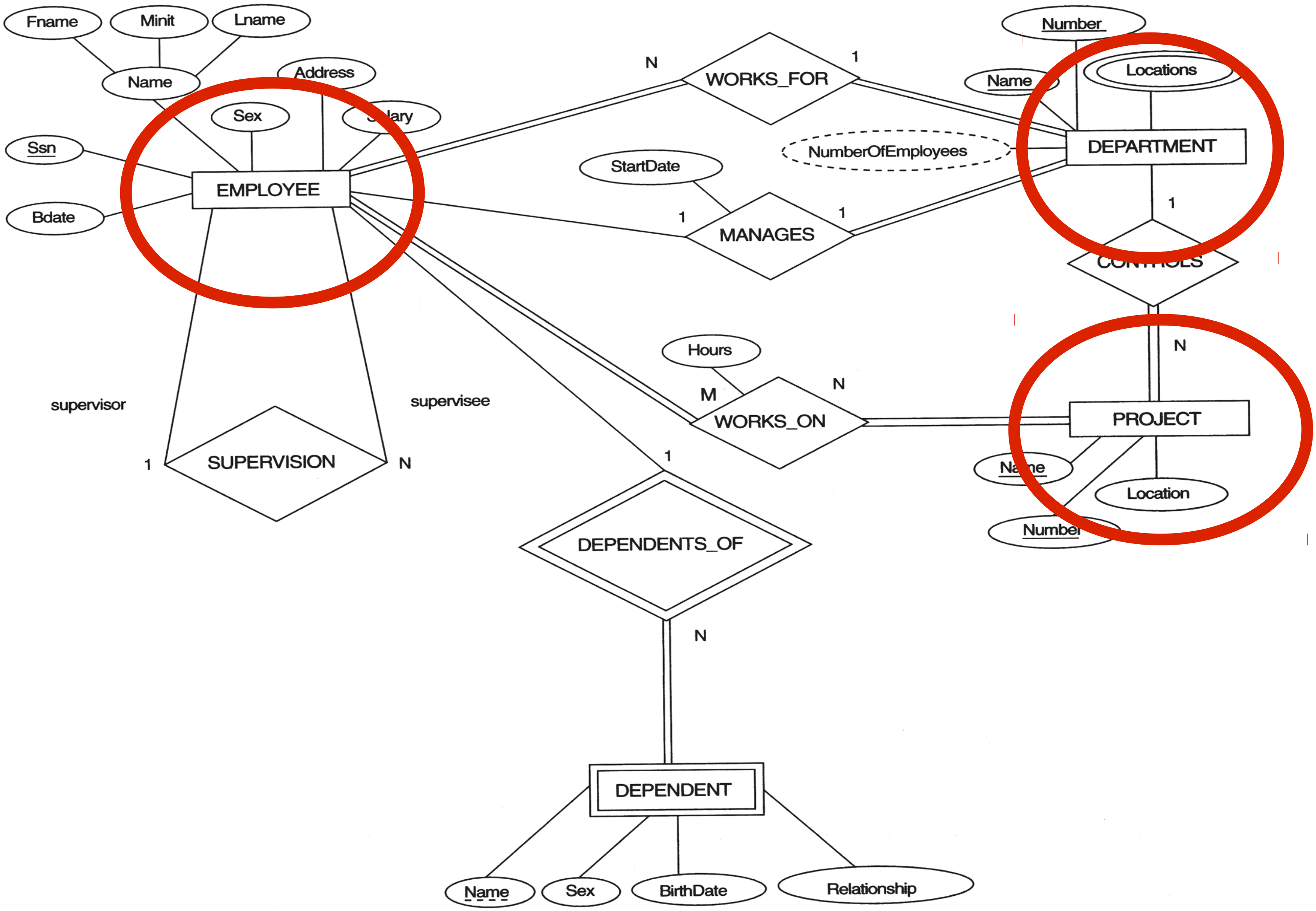
DEPENDENT

ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP

ER-to-Relational Mapping Algorithm

- **Step 1: Mapping of Regular Entity Types.**
  - For each regular (strong) entity type E in the ER schema, create a relation R that includes all the simple attributes of E.
  - Choose one of the key attributes of E as the primary key for R. If the chosen key of E is composite, the set of simple attributes that form it will together form the primary key of R.

- **Example:** We create the relations EMPLOYEE, DEPARTMENT, and PROJECT in the relational schema corresponding to the regular entities in the ER diagram. SSN, DNUMBER, and PNUMBER are the primary keys for the relations EMPLOYEE, DEPARTMENT, and PROJECT as shown.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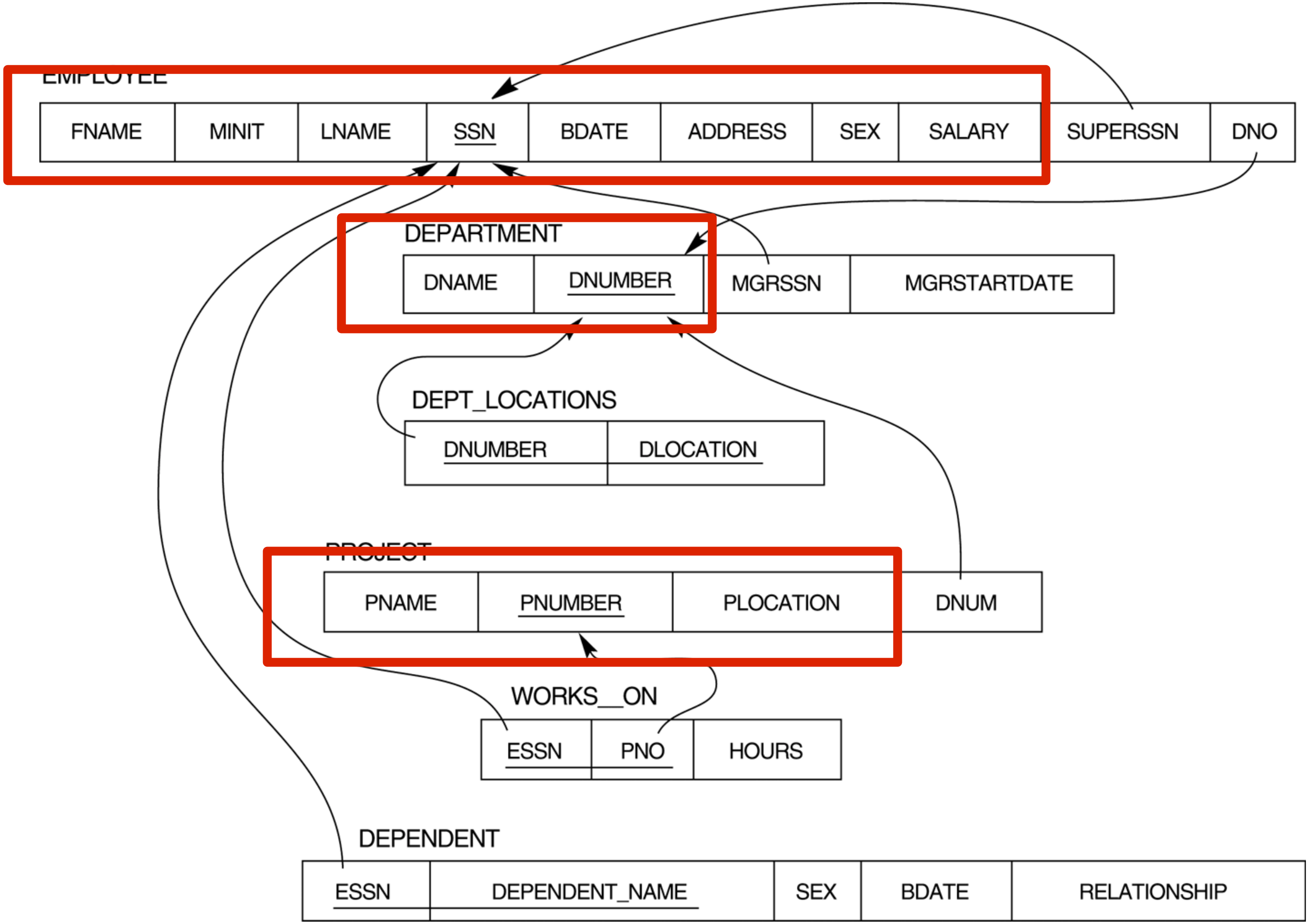


EMPLOYEE

FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO

DEPARTMENT

DNAME	DNUMBER	MGRSSN	MGRSTARTDATE

DEPT_LOCATIONS

DNUMBER	DLOCATION

PROJECT

PNAME	PNUMBER	PLOCATION	DNUM

WORKS__ON

ESSN	PNO	HOURS

DEPENDENT

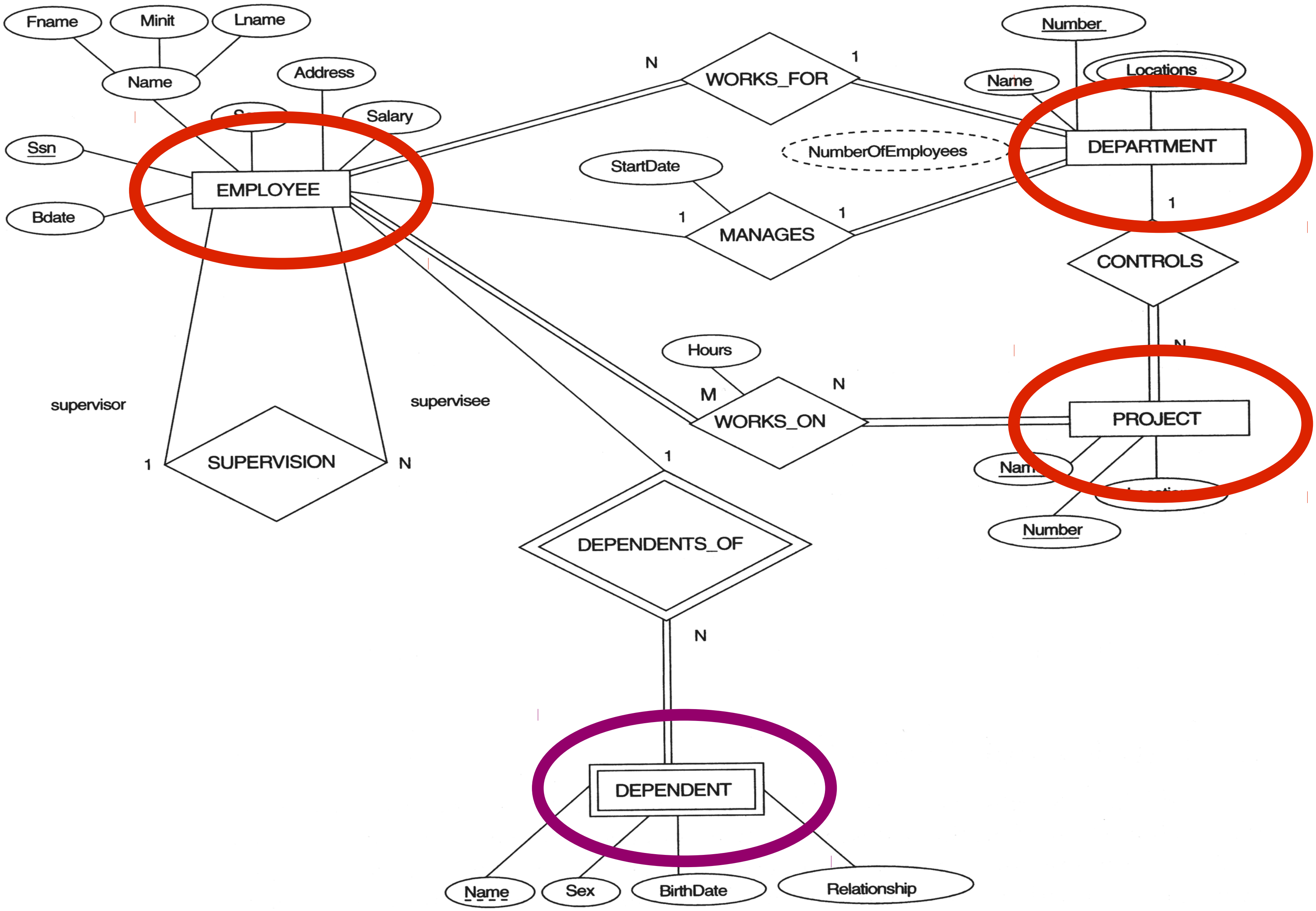
ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP

ER-to-Relational Mapping Algorithm (cont)

- **Step 2: Mapping of Weak Entity Types**
  - For each weak entity type W in the ER schema with owner entity type E, create a relation R and include all simple attributes (or simple components of composite attributes) of W as attributes of R.
  - In addition, include as foreign key attributes of R the primary key attribute(s) of the relation(s) that correspond to the owner entity type(s).
  - The primary key of R is the *combination of* the primary key(s) of the owner(s) and the partial key of the weak entity type W, if any.

- **Example:** Create the relation DEPENDENT in this step to correspond to the weak entity type DEPENDENT. Include the primary key SSN of the EMPLOYEE relation as a foreign key attribute of DEPENDENT (renamed to ESSN).

  The primary key of the DEPENDENT relation is the combination {ESSN, DEPENDENT_NAME} because DEPENDENT_NAME is the partial key of DEPENDENT.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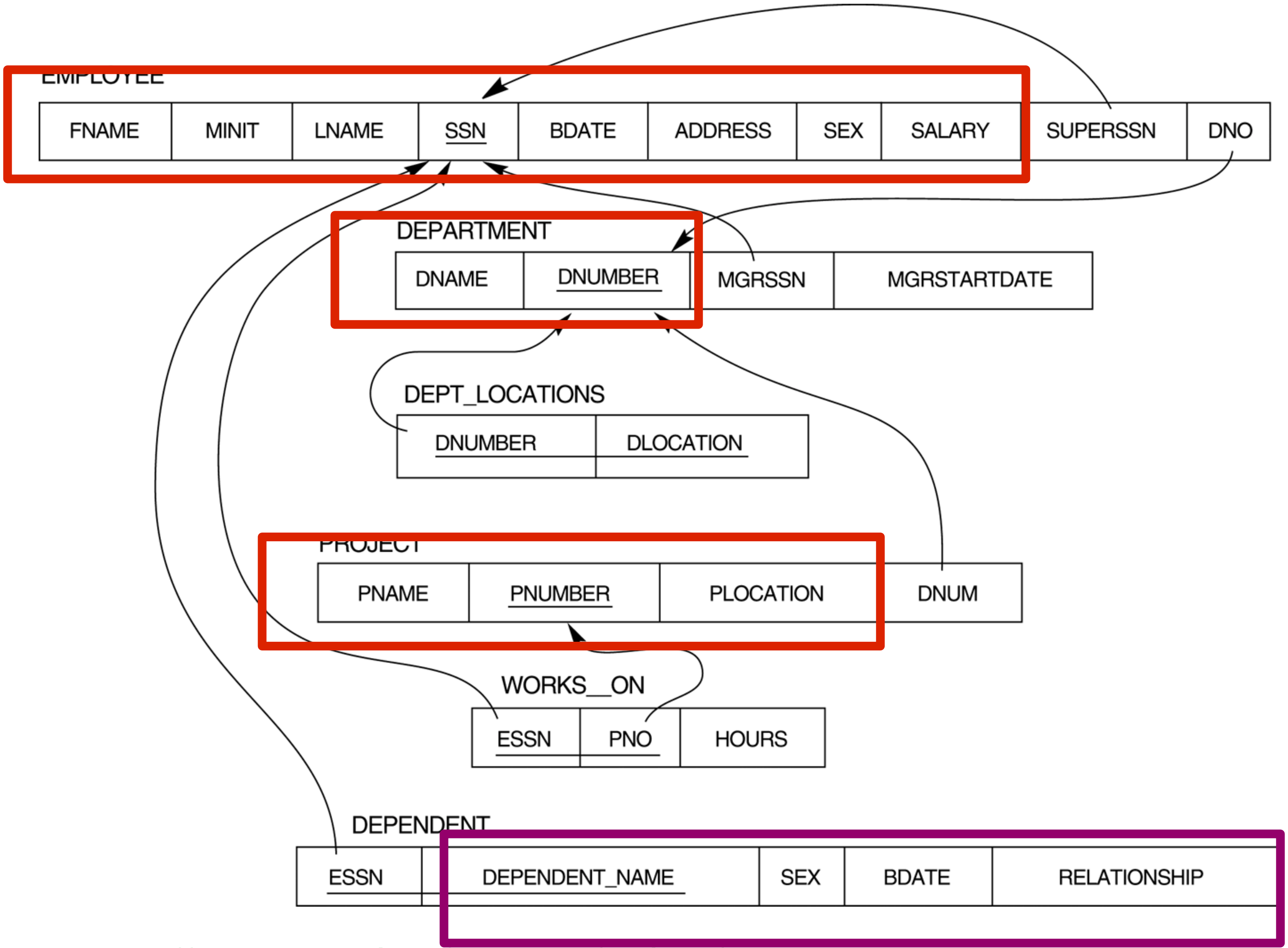


EMPLOYEE

FNAME MINIT LNAME SSN BDATE ADDRESS SEX SALARY SUPERSSN DNO

DEPARTMENT

DNAME DNUMBER MGRSSN MGRSTARTDATE

DEPT_LOCATIONS

DNUMBER DLOCATION

PROJECT

PNAME PNUMBER PLOCATION DNUM

WORKS__ON

ESSN PNO HOURS

DEPENDENT

ESSN DEPENDENT_NAME SEX BDATE RELATIONSHIP

ER-to-Relational Mapping Algorithm (cont)

- **Step 3: Mapping of Binary 1:1 Relation Types**

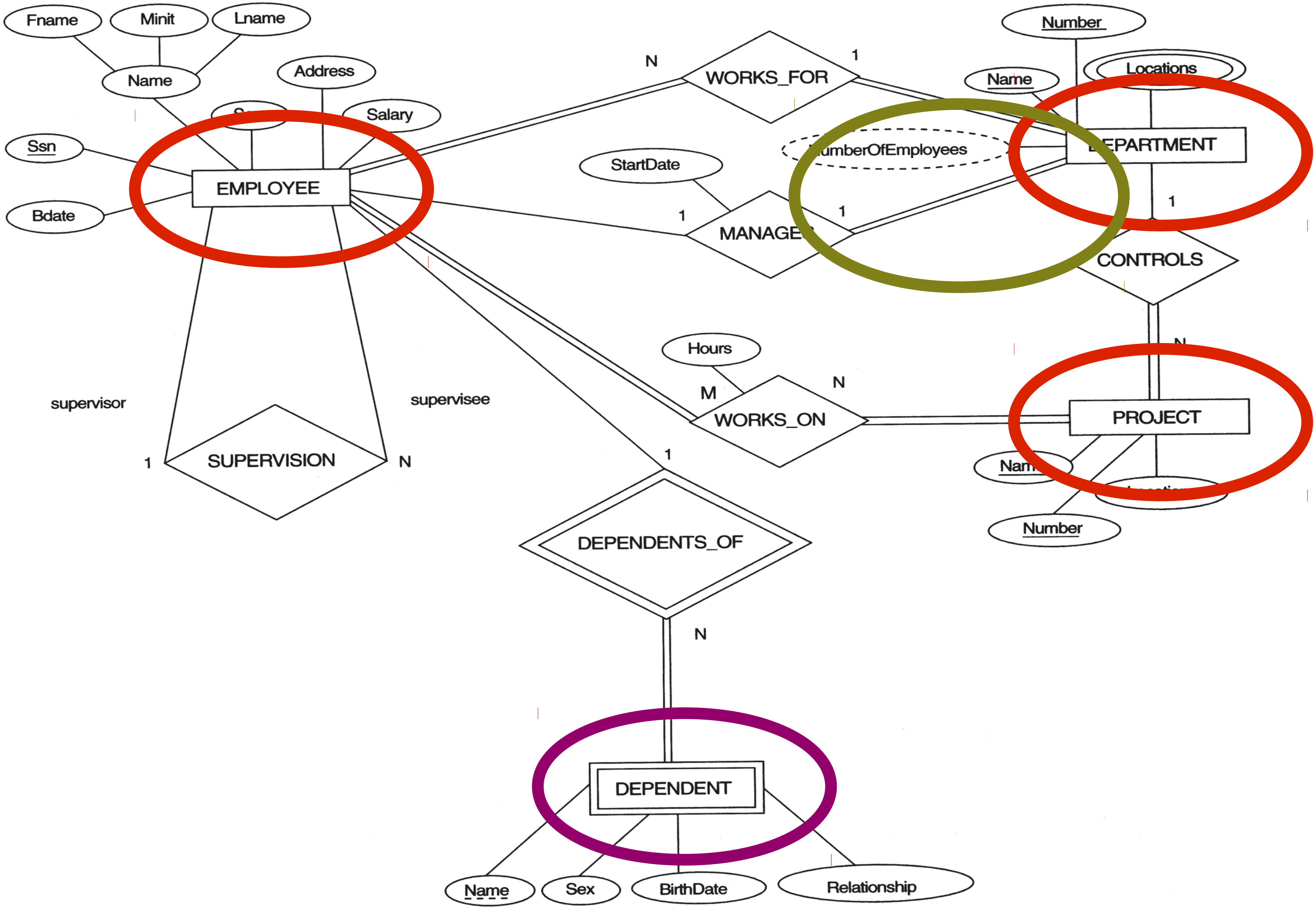
For each binary 1:1 relationship type R in the ER schema, identify the relations S and T that correspond to the entity types participating in R. There are three possible approaches:

(1) Foreign Key approach: Choose one of the relations-S, say-and include a foreign key in S the primary key of T. It is better to choose an entity type with *total participation* in R in the role of S.

Example: 1:1 relation MANAGES is mapped by choosing the participating entity type DEPARTMENT to serve in the role of S, because its participation in the MANAGES relationship type is total.

(2) Merged relation option: An alternate mapping of a 1:1 relationship type is possible by merging the two entity types and the relationship into a single relation. This may be appropriate when *both participations are total.*

(3) Cross-reference or relationship relation option: The third alternative is to set up a third relation R for the purpose of cross-referencing the primary keys of the two relations S and T representing the entity types.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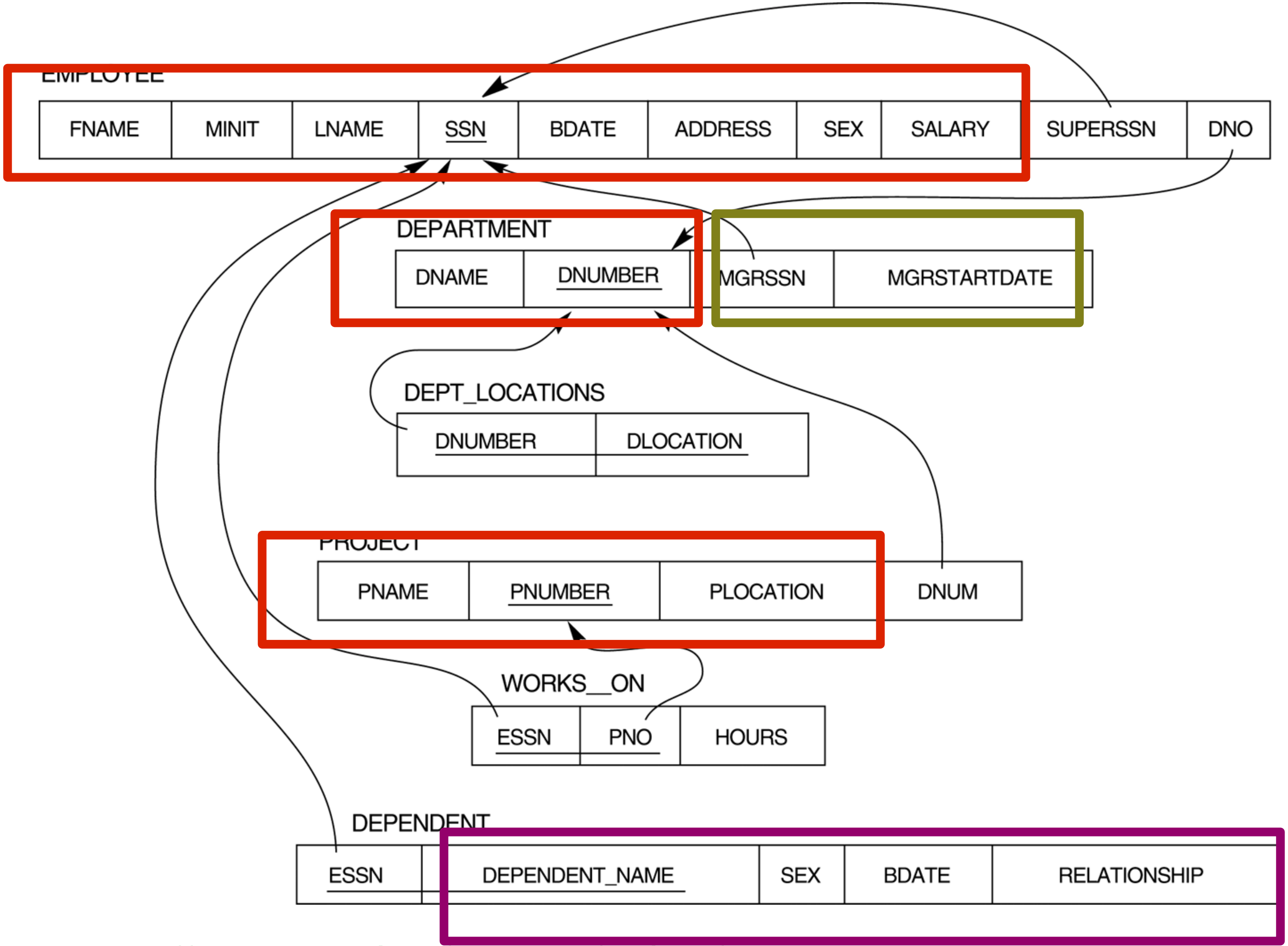


EMPLOYEE

FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO

DEPARTMENT

DNAME	DNUMBER	MGRSSN	MGRSTARTDATE

DEPT_LOCATIONS

DNUMBER	DLOCATION

PROJECT

PNAME	PNUMBER	PLOCATION	DNUM

WORKS__ON

ESSN	PNO	HOURS

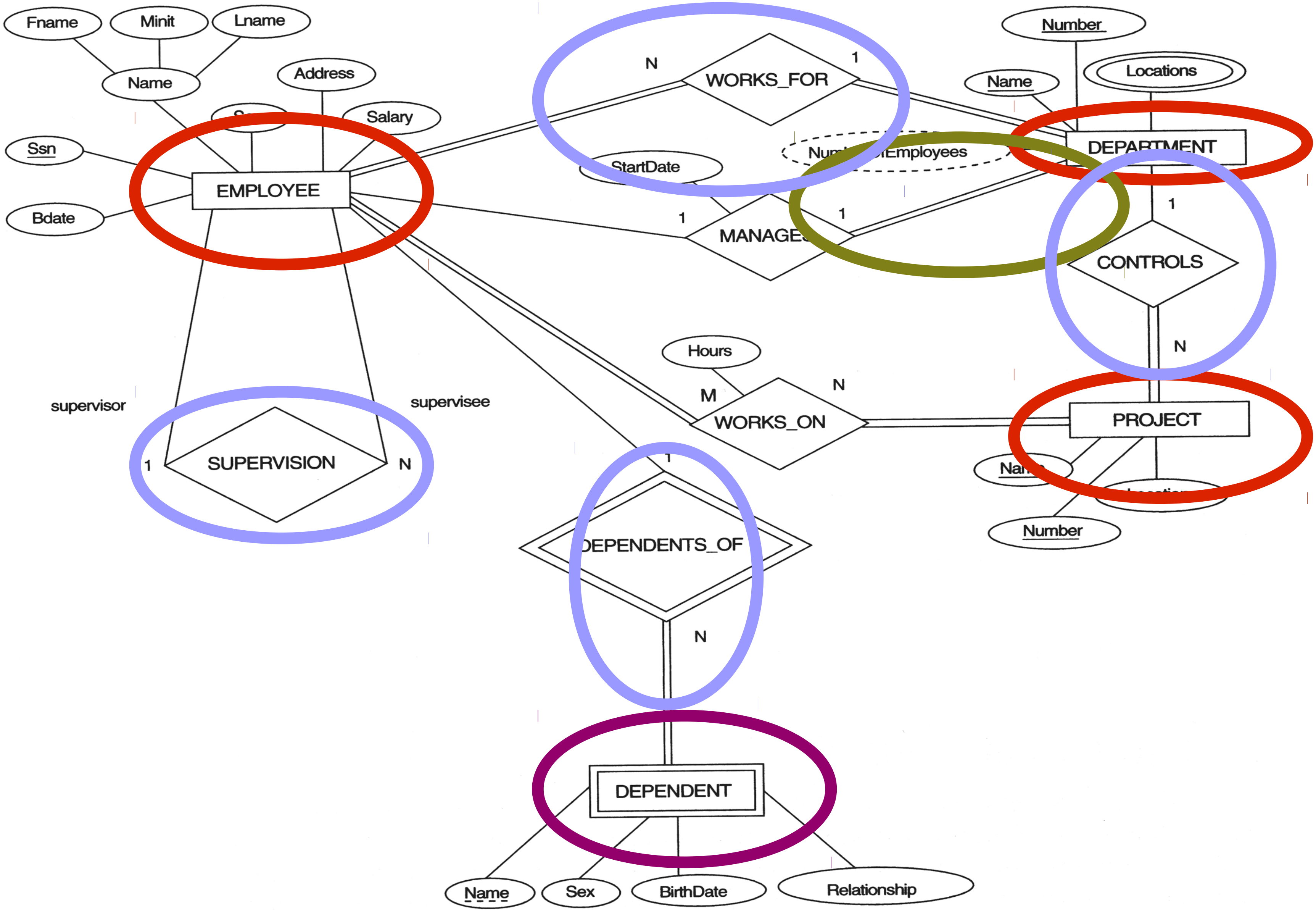
DEPENDENT

ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP

ER-to-Relational Mapping Algorithm (cont)

- **Step 4: Mapping of Binary 1:N Relationship Types.**
  - For each regular binary 1:N relationship type R, identify the relation S that represent the participating entity type at the N-side of the relationship type.
  - Include as foreign key in S the primary key of the relation T that represents the other entity type participating in R.
  - Include any simple attributes of the 1:N relation type as attributes of S.

- Example: 1:N relationship types WORKS_FOR, CONTROLS, and SUPERVISION in the figure. For WORKS_FOR we include the primary key DNUMBER of the DEPARTMENT relation as foreign key in the EMPLOYEE relation and call it DNO.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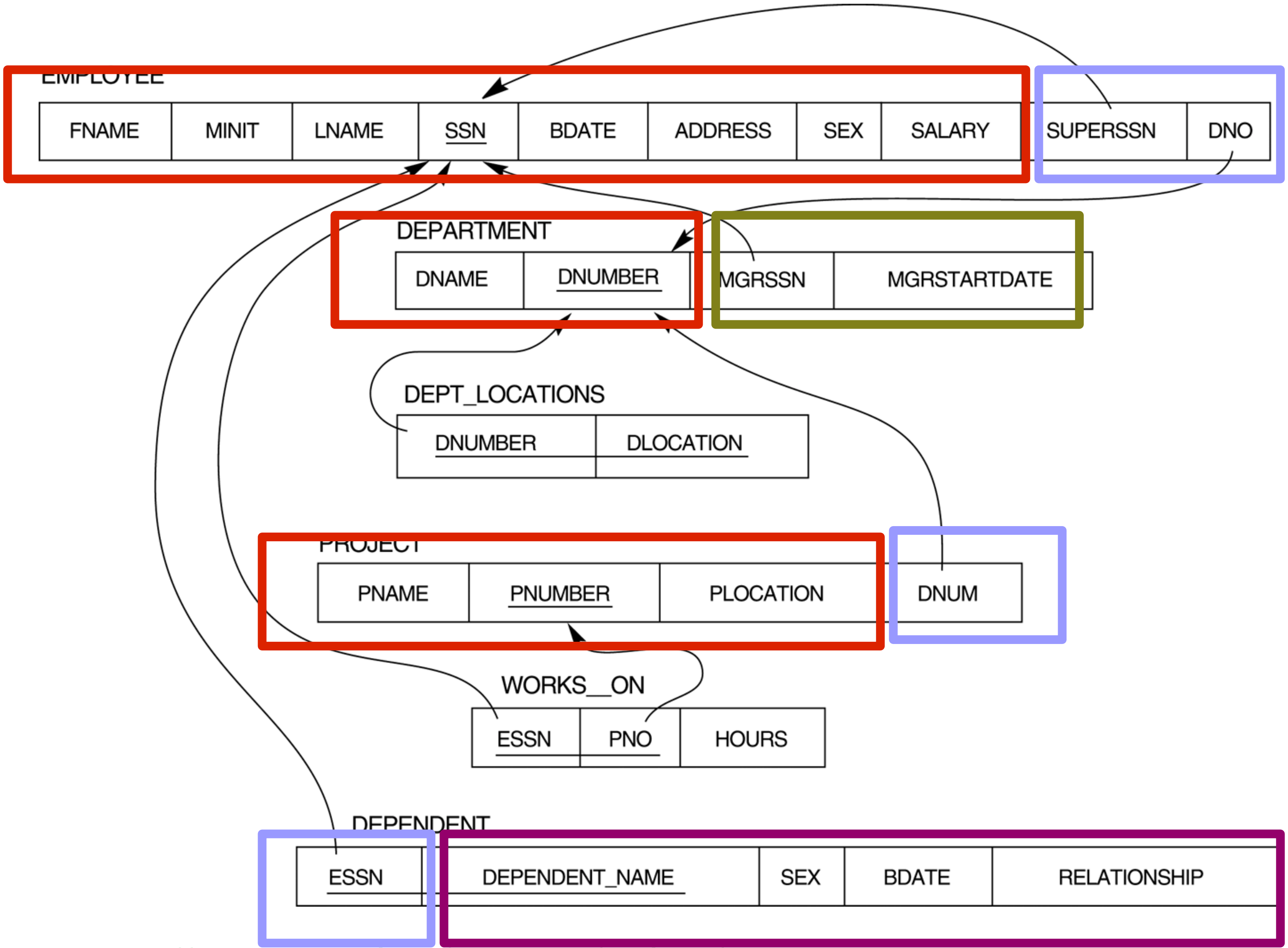


EMPLOYEE

FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO

DEPARTMENT

DNAME	DNUMBER	MGRSSN	MGRSTARTDATE

DEPT_LOCATIONS

DNUMBER	DLOCATION

PROJECT

PNAME	PNUMBER	PLOCATION	DNUM

WORKS__ON

ESSN	PNO	HOURS

DEPENDENT

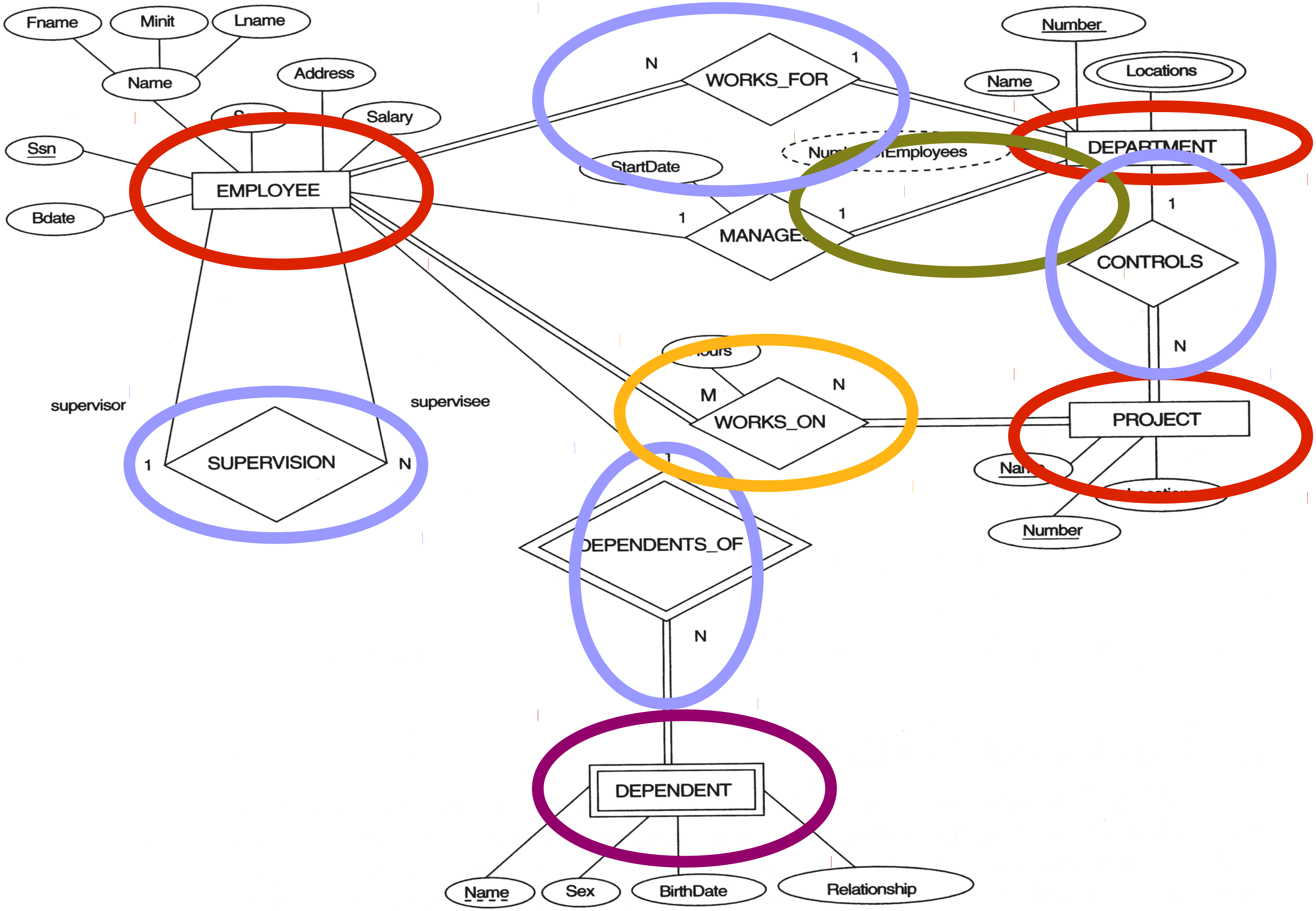
ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP

ER-to-Relational Mapping Algorithm (cont)

- **Step 5: Mapping of Binary M:N Relationship Types.**
  - For each regular binary M:N relationship type R, *create a new relation* S to represent R.
  - Include as foreign key attributes in S the primary keys of the relations that represent the participating entity types; *their combination will form the primary key* of S.
  - Also include any simple attributes of the M:N relationship type (or simple components of composite attributes) as attributes of S.

- **Example:** The M:N relationship type WORKS_ON from the ER diagram is mapped by creating a relation WORKS_ON in the relational database schema. The primary keys of the PROJECT and EMPLOYEE relations are included as foreign keys in WORKS_ON and renamed PNO and ESSN, respectively.

  Attribute HOURS in WORKS_ON represents the HOURS attribute of the relation type. The primary key of the WORKS_ON relation is the combination of the foreign key attributes {ESSN, PNO}.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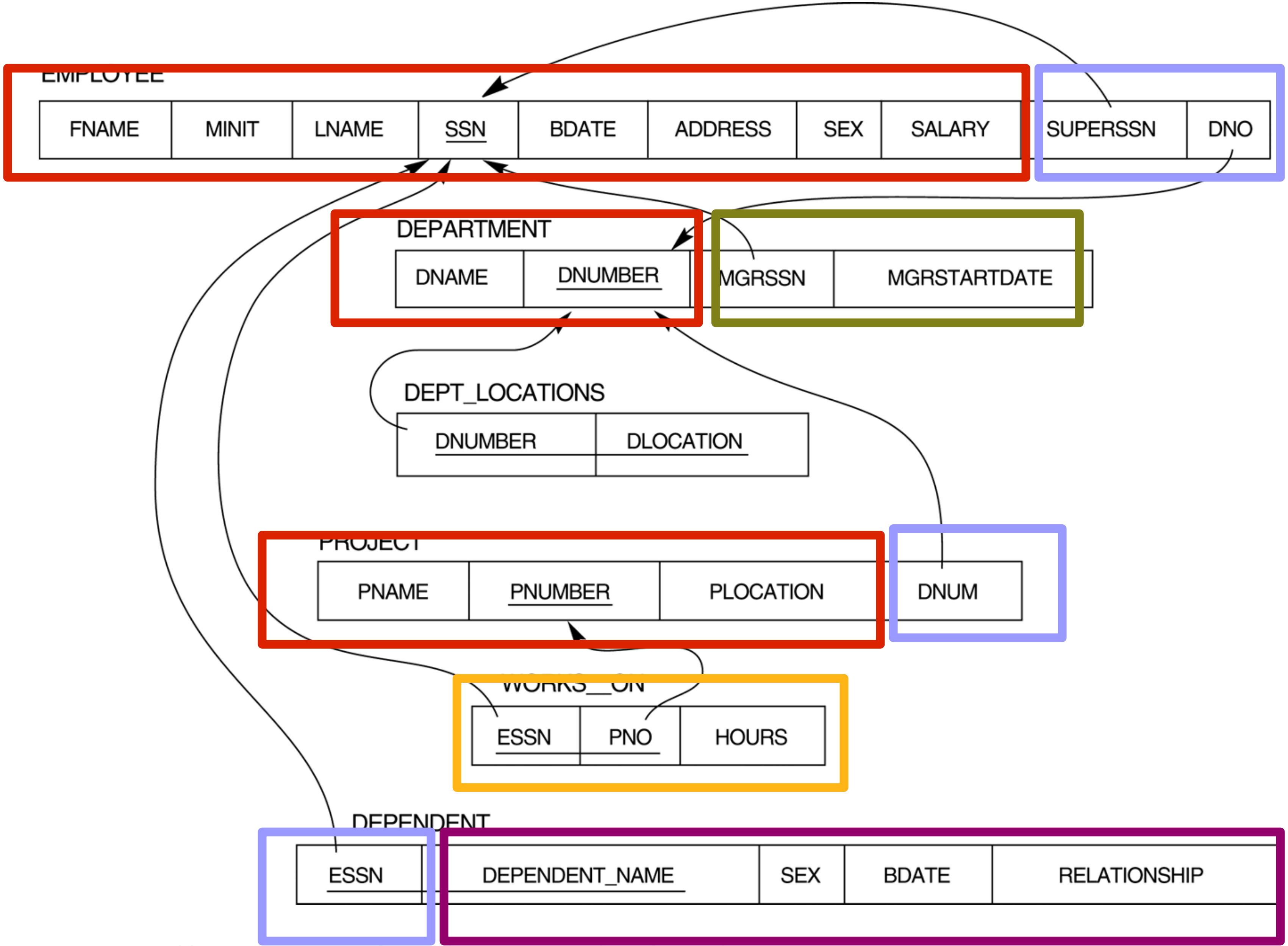


EMPLOYEE

FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO

DEPARTMENT

DNAME	DNUMBER	MGRSSN	MGRSTARTDATE

DEPT_LOCATIONS

DNUMBER	DLOCATION

PROJECT

PNAME	PNUMBER	PLOCATION	DNUM

WORKS_ON

ESSN	PNO	HOURS

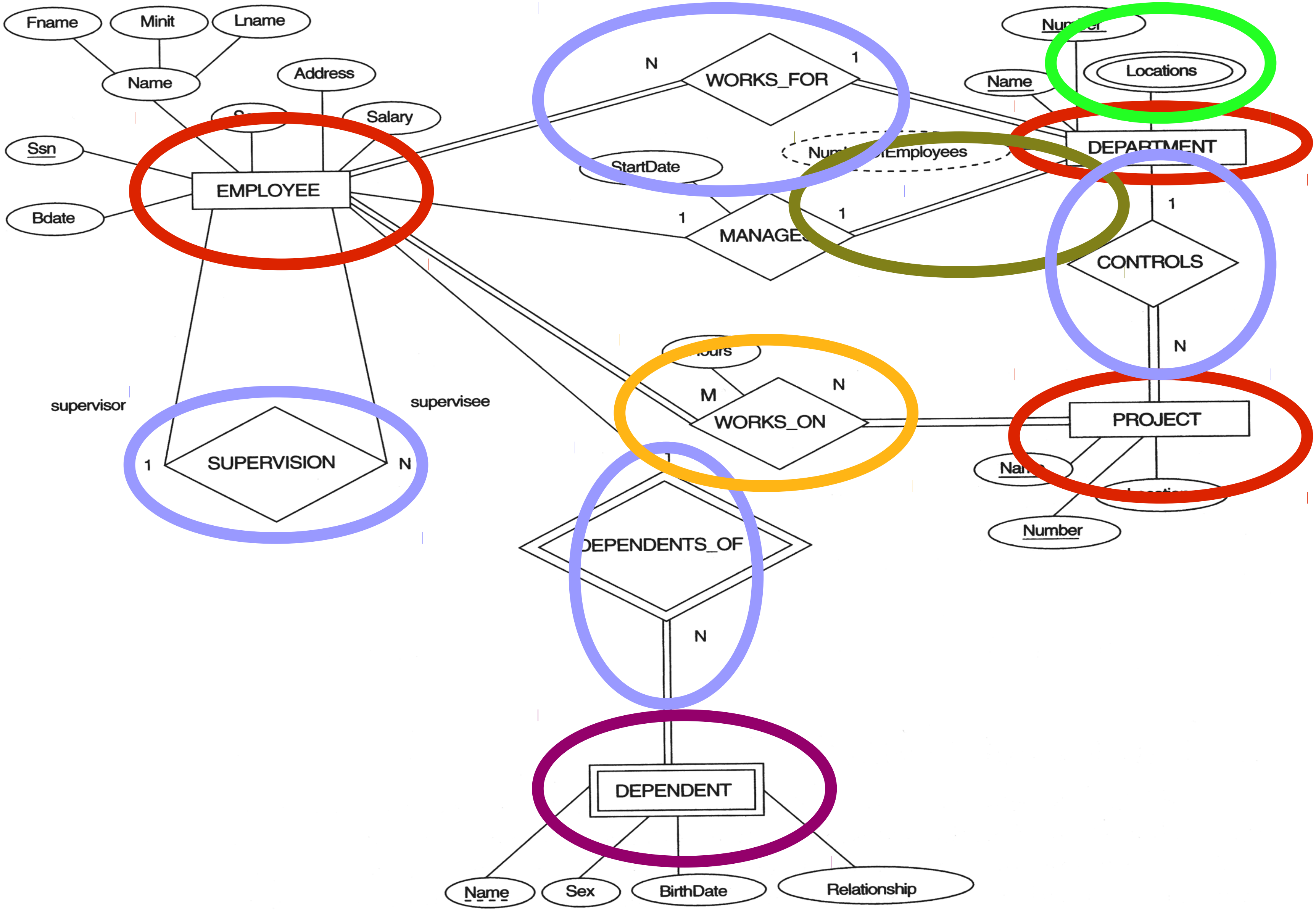
DEPENDENT

ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP

ER-to-Relational Mapping Algorithm (cont)

- **Step 6: Mapping of Multivalued attributes.**
  - For each multivalued attribute A, create a new relation R. This relation R will include an attribute corresponding to A, plus the primary key attribute K-as a foreign key in R-of the relation that represents the entity type of relationship type that has A as an attribute.
  - The primary key of R is the combination of A and K. If the multivalued attribute is composite, we include its simple components.

- **Example:** The relation DEPT_LOCATIONS is created. The attribute DLOCATION represents the multivalued attribute LOCATIONS of DEPARTMENT, while DNUMBER-as foreign key-represents the primary key of the DEPARTMENT relation. The primary key of R is the combination of {DNUMBER, DLOCATION}.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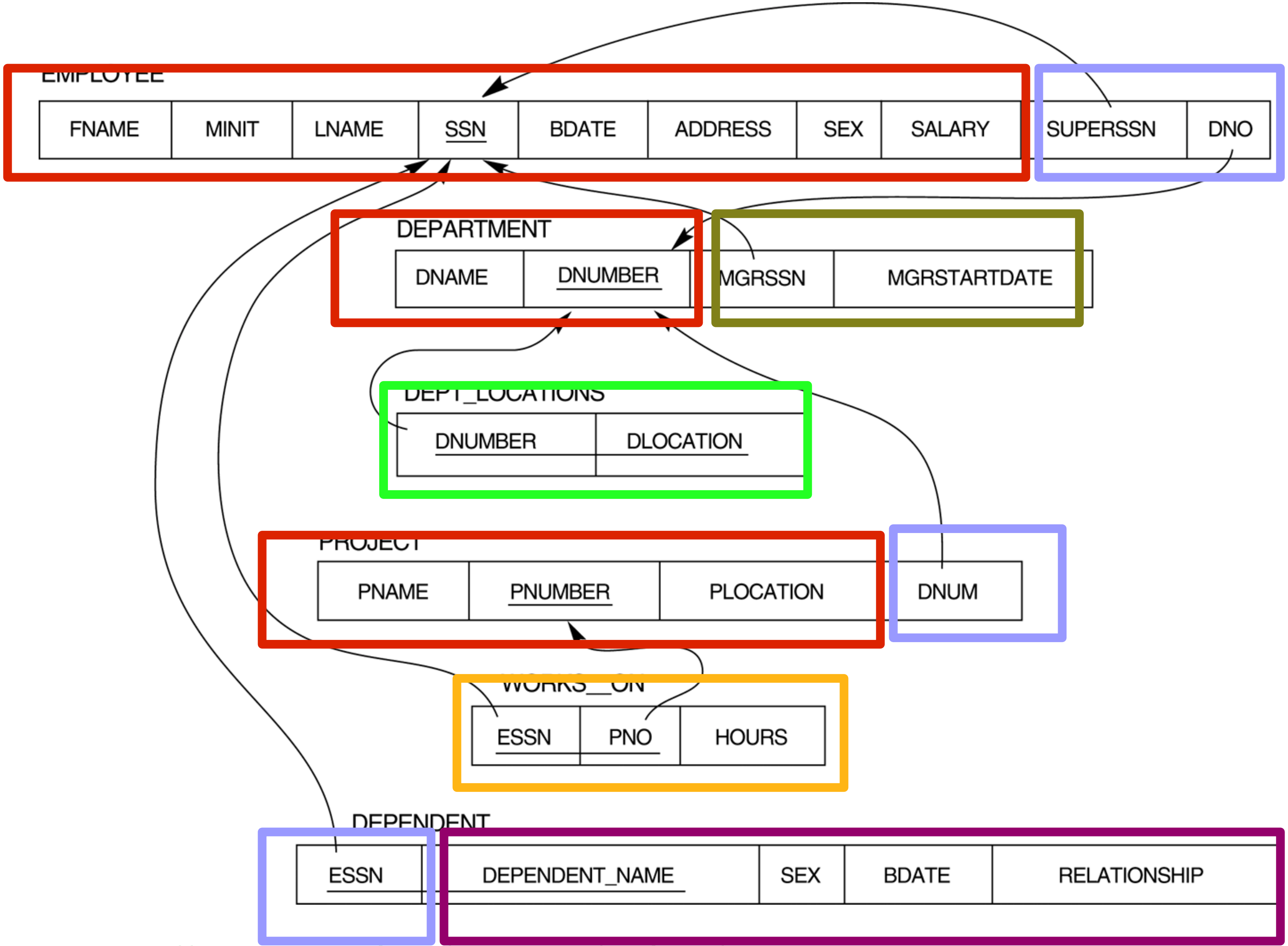


EMPLOYEE

FNAME	MINIT	LNAME	SSN	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO

DEPARTMENT

DNAME	DNUMBER	MGRSSN	MGRSTARTDATE

DEPT_LOCATIONS

DNUMBER	DLOCATION

PROJECT

PNAME	PNUMBER	PLOCATION	DNUM

WORKS_ON

ESSN	PNO	HOURS

DEPENDENT

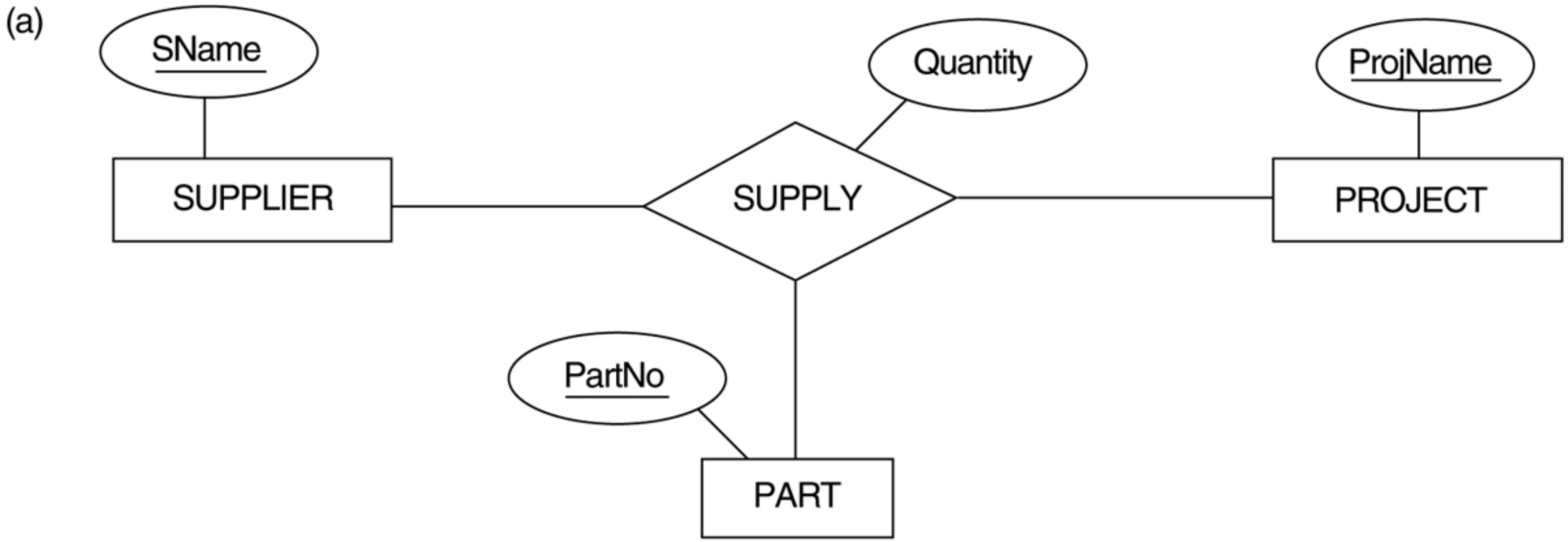
ESSN	DEPENDENT_NAME	SEX	BDATE	RELATIONSHIP

ER-to-Relational Mapping Algorithm (cont)

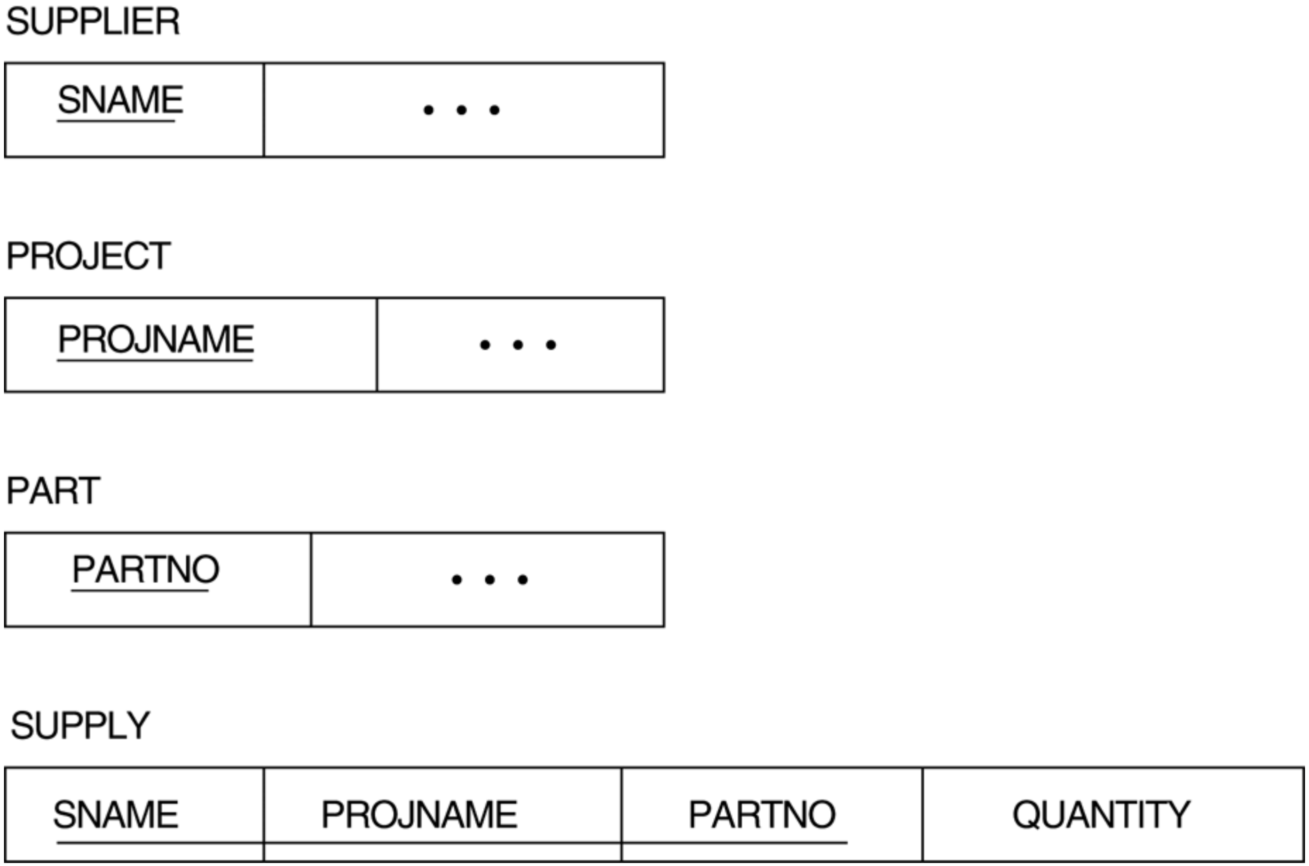
- **Step 7: Mapping of N-ary Relationship Types.**
  - For each n-ary relationship type R, where n>2, create a new relationship S to represent R.
  - Include as foreign key attributes in S the primary keys of the relations that represent the participating entity types.
  - Also include any simple attributes of the n-ary relationship type (or simple components of composite attributes) as attributes of S.

- **Example:** The relationship type SUPPY in the ER below. This can be mapped to the relation SUPPLY shown in the relational schema, whose primary key is the combination of the three foreign keys {SNAME, PARTNO, PROJNAME}

Ternary relationship types. (a) The SUPPLY relationship.



Mapping the *n*-ary relationship type SUPPLY from prev diagram



SUPPLIER

SNAME	. . .

PROJECT

PROJNAME	. . .

PART

PARTNO	. . .

SUPPLY

SNAME	PROJNAME	PARTNO	QUANTITY

Summary of Mapping constructs and constraints

Correspondence between ER and Relational Models

ER Model	Relational Model
Entity type	“Entity” relation
1:1 or 1:N relationship type	Foreign key (or “relationship” relation)
M:N relationship type	“Relationship” relation and two foreign keys
n-ary relationship type	“Relationship” relation and n foreign keys
Simple attribute	Attribute
Composite attribute	Set of simple component attributes
Multivalued attribute	Relation and foreign key
Value set	Domain
Key attribute	Primary (or secondary) key

Mapping EER Model Constructs to Relations

- **Step8: Options for Mapping Specialization or Generalization.**

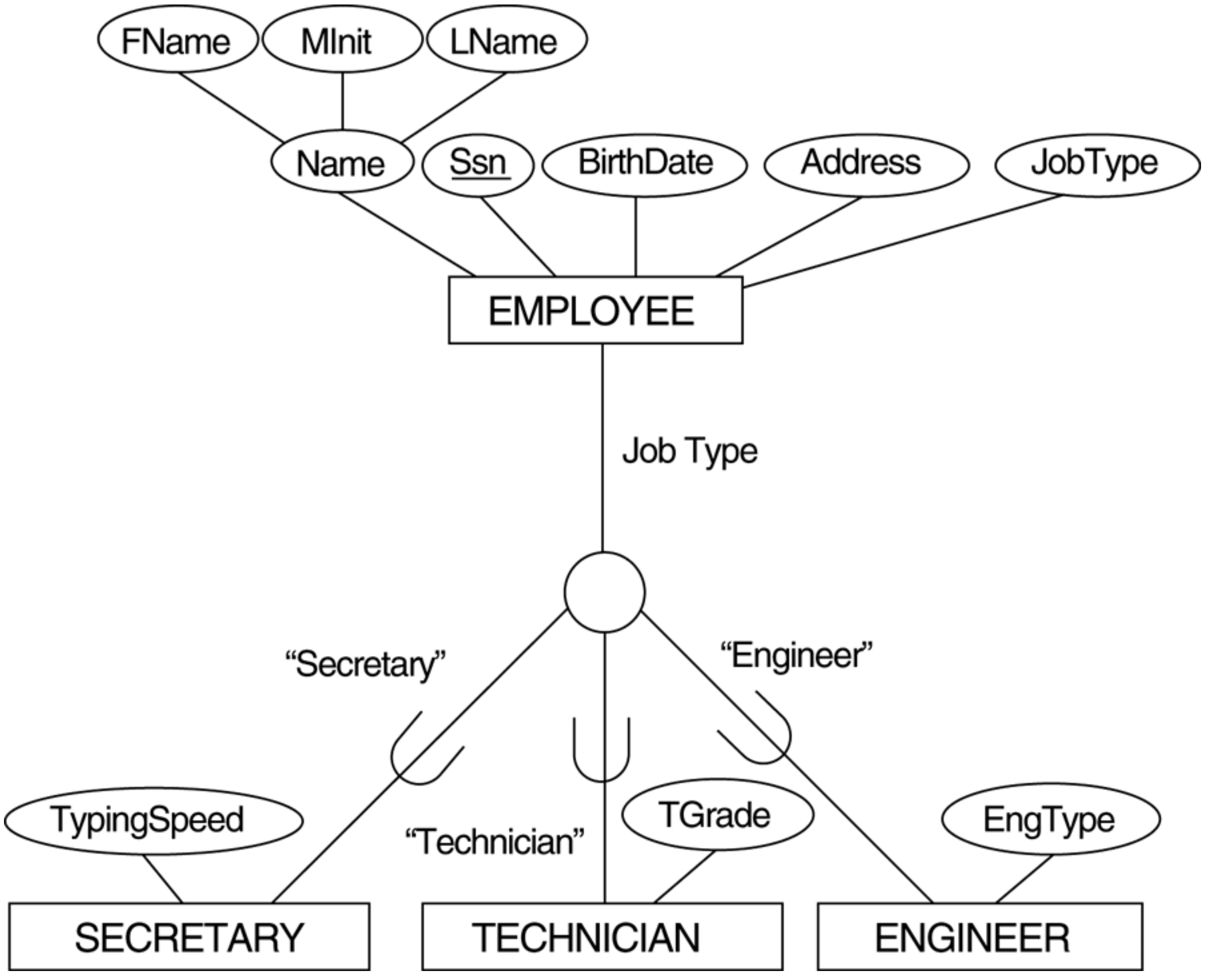
  Convert each specialization with m subclasses $\{S_1, S_2, \ldots, S_m\}$ and generalized superclass C, where the attributes of C are $\{k, a_1, \ldots a_n\}$ and k is the (primary) key, into relational schemas using one of the four following options:

- Option 8A: Multiple relations-Superclass and subclasses. Create a relation L for C with attributes Attrs(L) = {k,a_1,…a_n} and PK(L) = k. Create a relation L_i for each subclass S_i, 1 < i < m, with the attributesAttrs(L_i) = {k} U {attributes of S_i} and PK(L_i)=k. This option works **for any specialization** (total or partial, disjoint of over-lapping).

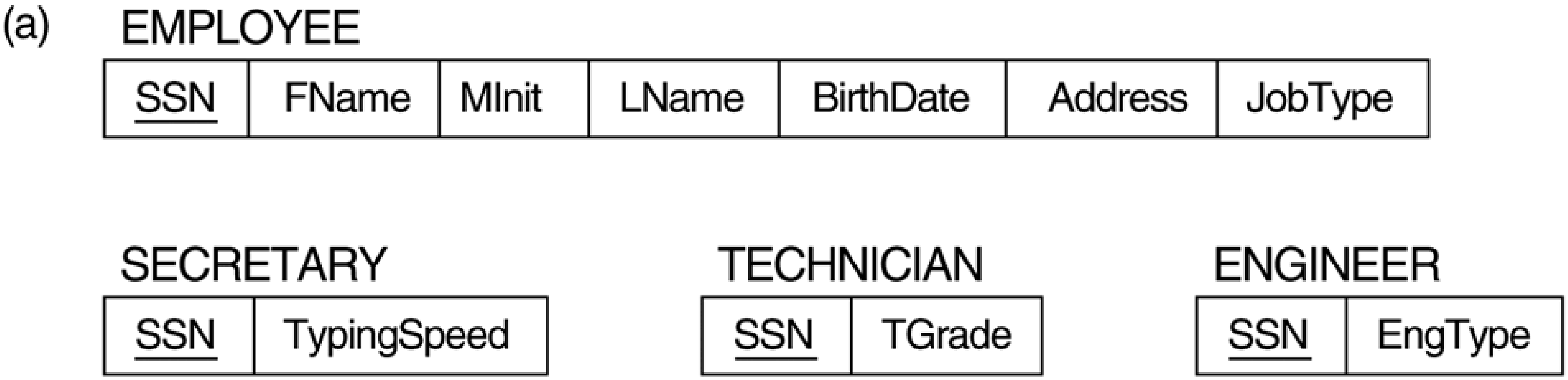
- Option 8B: Multiple relations-Subclass relations only

  Create a relation L_i for each subclass S_i, 1 < i < m, with the attributes Attr(L_i) = {attributes of S_i} U {k,a_1…,a_n} and PK(L_i) = k. This option only works for a specialization whose subclasses are **total** (every entity in the superclass must belong to (at least) one of the subclasses).

EER diagram notation for an attribute-defined specialization on JobType.



Options for mapping specialization or generalization.
(a) Mapping the EER schema using option 8A.



(a)

EMPLOYEE

SSN	FName	MInit	LName	BirthDate	Address	JobType

SECRETARY

SSN	TypingSpeed

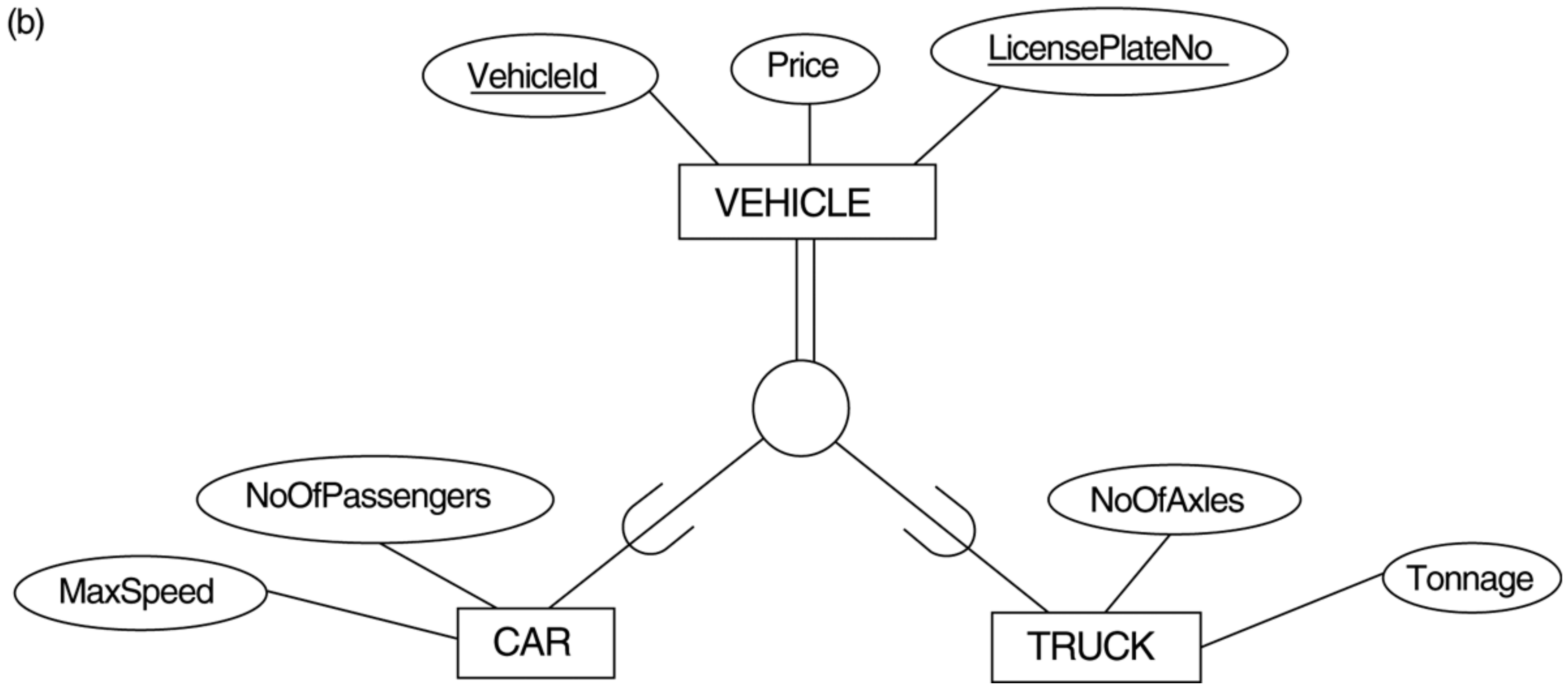
TECHNICIAN

SSN	TGrade

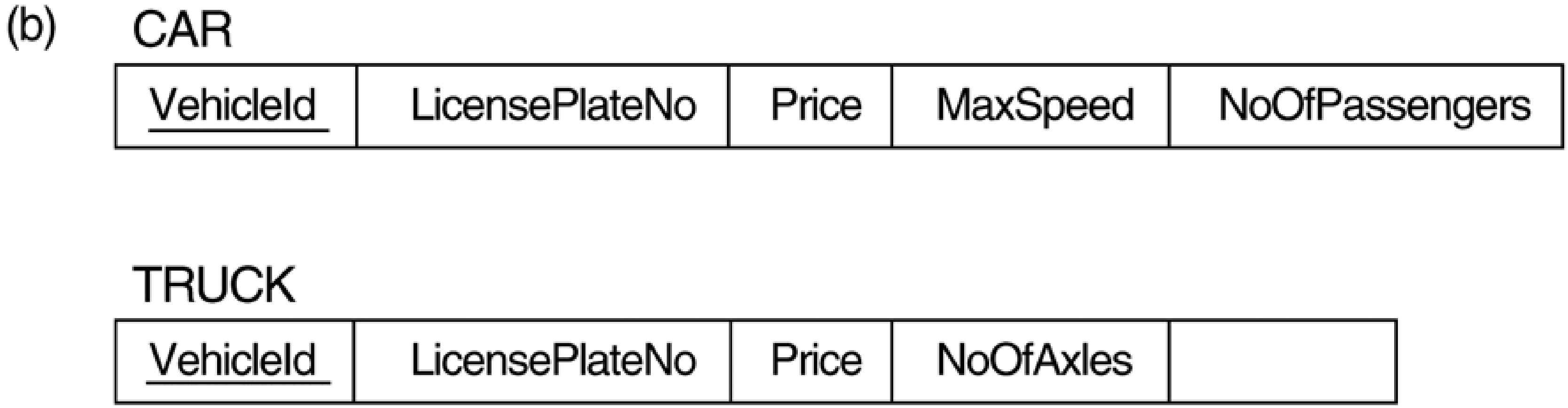
ENGINEER

SSN	EngType

Generalization. (b) Generalizing CAR and TRUCK into the superclass VEHICLE.



Options for mapping specialization or generalization. (b) Mapping the EER schema using option 8B.



(b)

CAR

VehicleId	LicensePlateNo	Price	MaxSpeed	NoOfPassengers

TRUCK

VehicleId	LicensePlateNo	Price	NoOfAxles	

Mapping EER Model Constructs to Relations (cont)

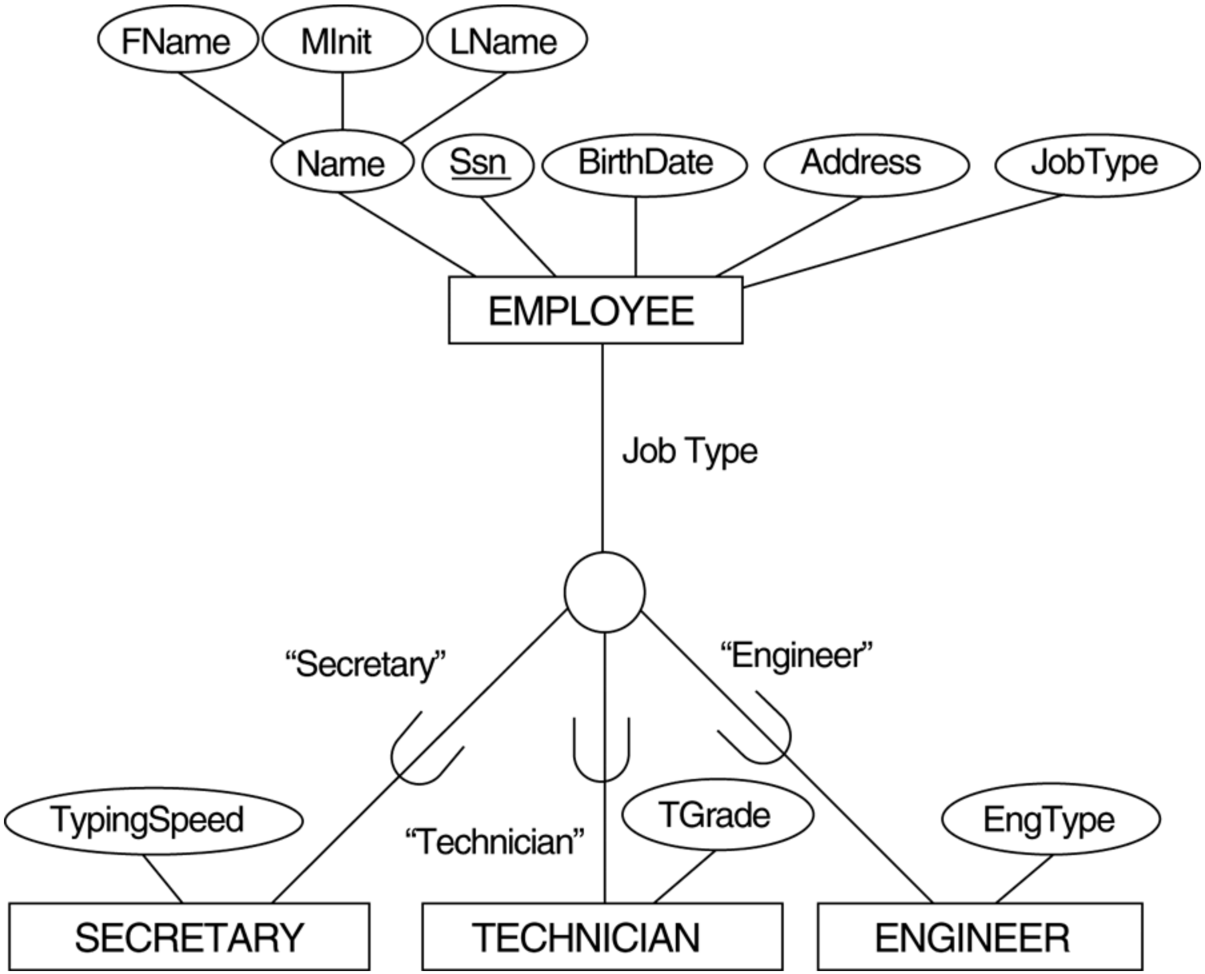
Option 8C: Single relation with one type attribute.

Create a single relation L with attributes Attrs(L) = $\{k,a_1,\ldots a_n\}$ U {attributes of S_1} U…U {attributes of S_m} U {t} and PK(L) = k. The attribute t is called a type (or **discriminating**) attribute that indicates the subclass to which each tuple belongs

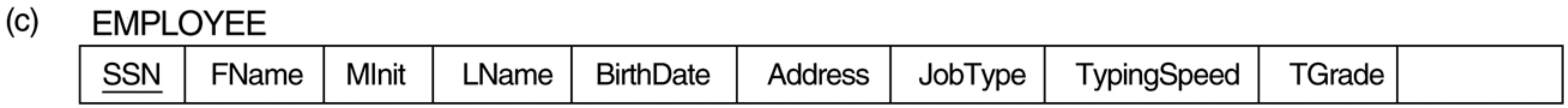
Option 8D: Single relation with multiple type attributes.

Create a single relation schema L with attributes Attrs(L) = $\{k,a_1,\ldots a_n\}$ U {attributes of S_1} U…U {attributes of S_m} U $\{t_1, t_2,\ldots,t_m\}$ and PK(L) = k. Each t_i, 1 < I < m, is a Boolean type attribute indicating whether a tuple belongs to the subclass S_i.

EER diagram notation for an attribute-defined specialization on JobType.



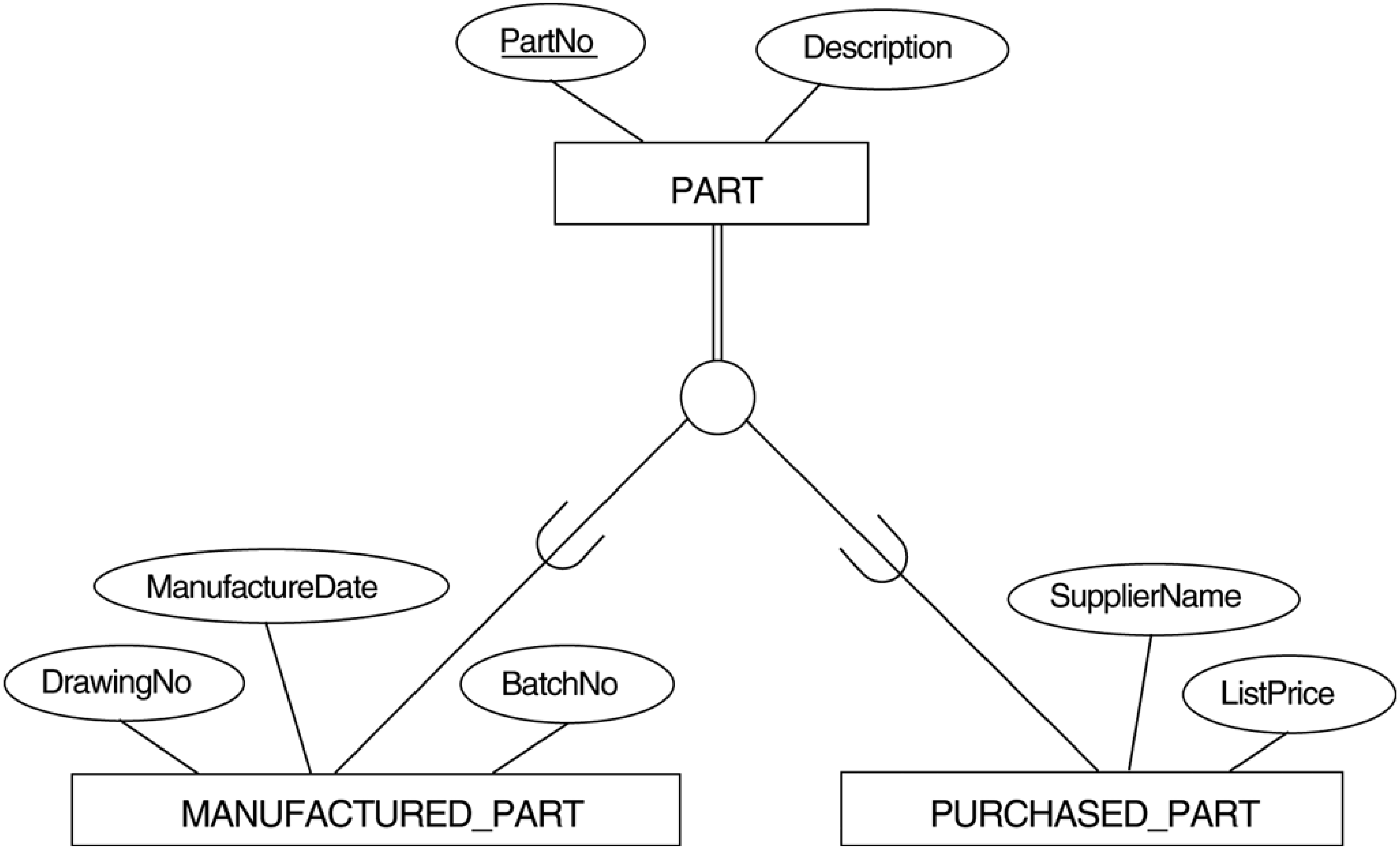
Options for mapping specialization or generalization. (c) Mapping the EER schema using option 8C.



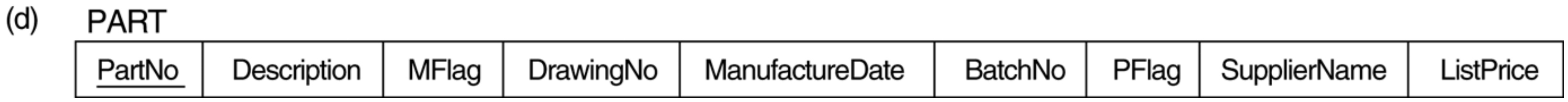
(c) EMPLOYEE

SSN	FName	MInit	LName	BirthDate	Address	JobType	TypingSpeed	TGrade	

EER diagram notation for an overlapping (nondisjoint) specialization.



Options for mapping specialization or generalization. (d) Mapping using option 8D with Boolean type fields Mflag and Pflag.



(d) PART

PartNo	Description	MFlag	DrawingNo	ManufactureDate	BatchNo	PFlag	SupplierName	ListPrice

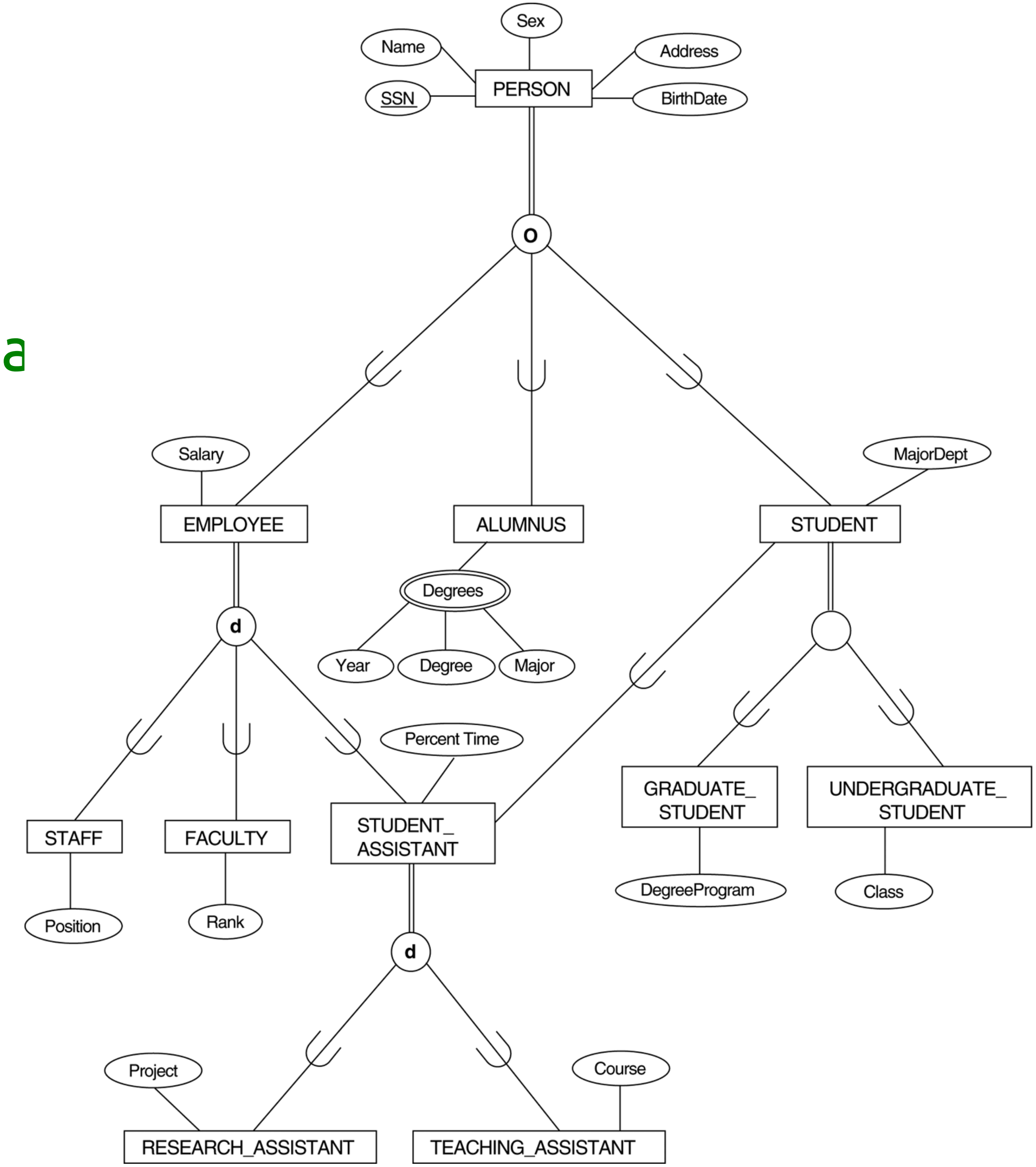
Mapping EER Model Constructs to Relations (cont)

- **Mapping of Shared Subclasses (Multiple Inheritance)**

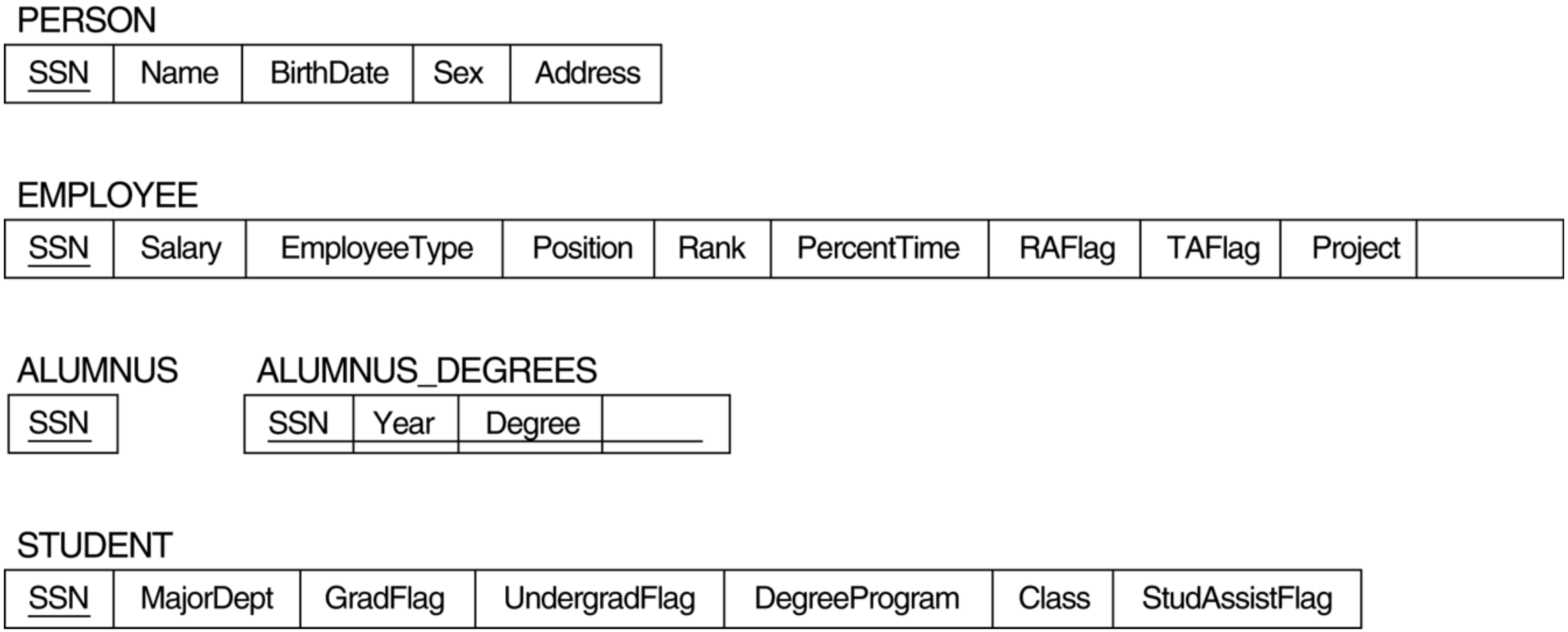
  A shared subclass, such as STUDENT_ASSISTANT, is a subclass of several classes, indicating multiple inheritance. These classes must all have the same key attribute; otherwise, the shared subclass would be modeled as a category.

  We can apply any of the options discussed in Step 8 to a shared subclass, subject to the restriction discussed in Step 8 of the mapping algorithm. Below both 8C and 8D are used for the shared class STUDENT_ASSISTANT.

A specialization lattice with multiple inheritance for a UNIVERSITY database.



Mapping the EER specialization lattice in prev diagram using multiple options.



PERSON

SSN	Name	BirthDate	Sex	Address

EMPLOYEE

SSN	Salary	EmployeeType	Position	Rank	PercentTime	RAFlag	TAFlag	Project	

ALUMNUS

SSN

ALUMNUS_DEGREES

SSN	Year	Degree	

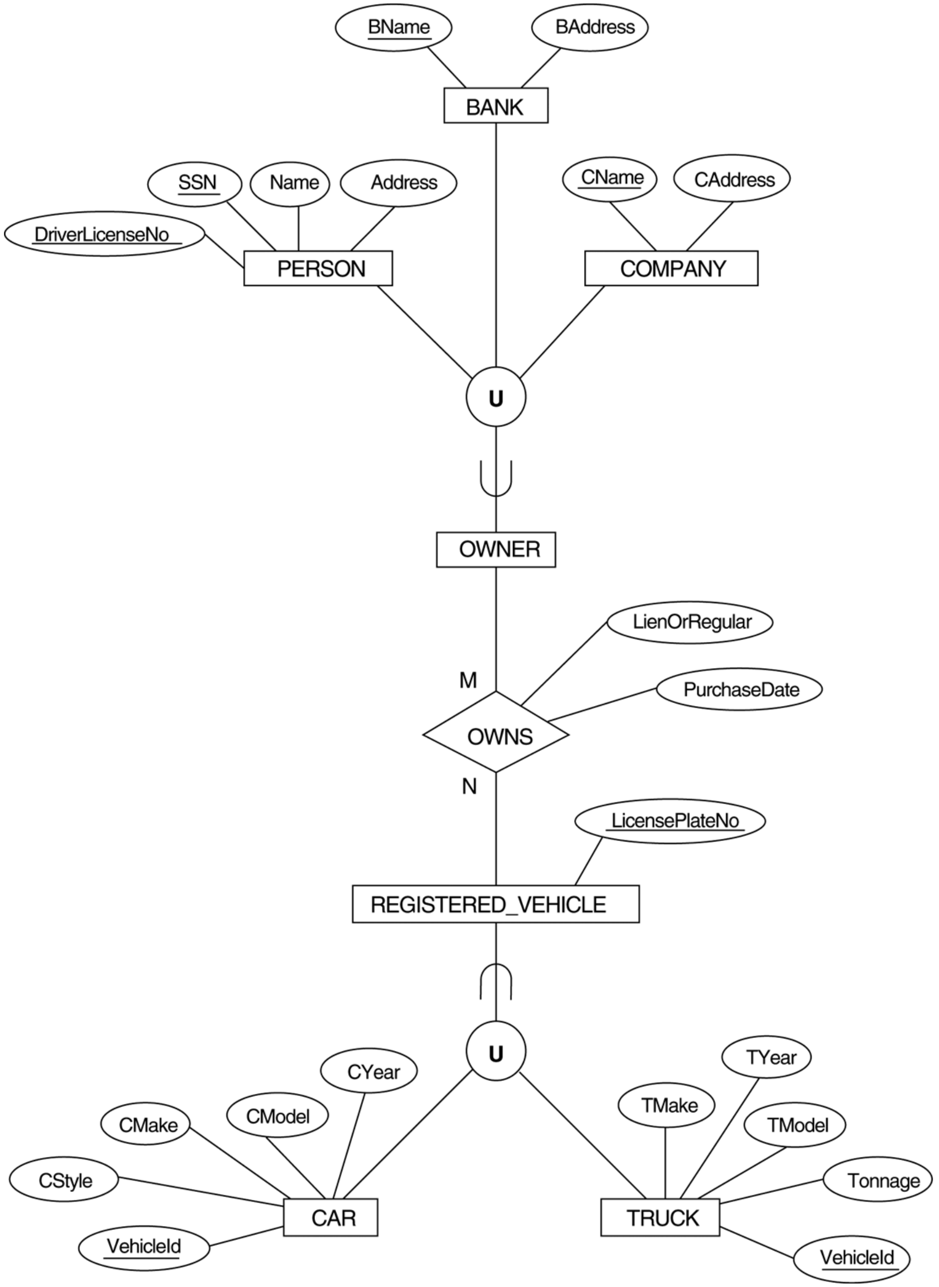
STUDENT

SSN	MajorDept	GradFlag	UndergradFlag	DegreeProgram	Class	StudAssistFlag

Mapping EER Model Constructs to Relations (cont)

- **Step 9: Mapping of Union Types (Categories).**
  - For mapping a category whose defining superclass have different keys, it is customary to specify a new key attribute, called a **surrogate key**, when creating a relation to correspond to the category.
  - In the example below we can create a relation OWNER to correspond to the OWNER category and include any attributes of the category in this relation. The primary key of the OWNER relation is the surrogate key, which we called OwnerId.

Two categories (union types): OWNER and REGISTERED_VEHICLE.



Mapping the EER categories (union types) to relations.

PERSON

SSN	DriverLicenseNo	Name	Address	

BANK

BName	BAddress	OwnerId

COMPANY

CName	CAddress	OwnerId

OWNER

OwnerId

REGISTERED_VEHICLE

VehicleId	LicensePlateNumber

CAR

VehicleId	CStyle	CMake	CModel	

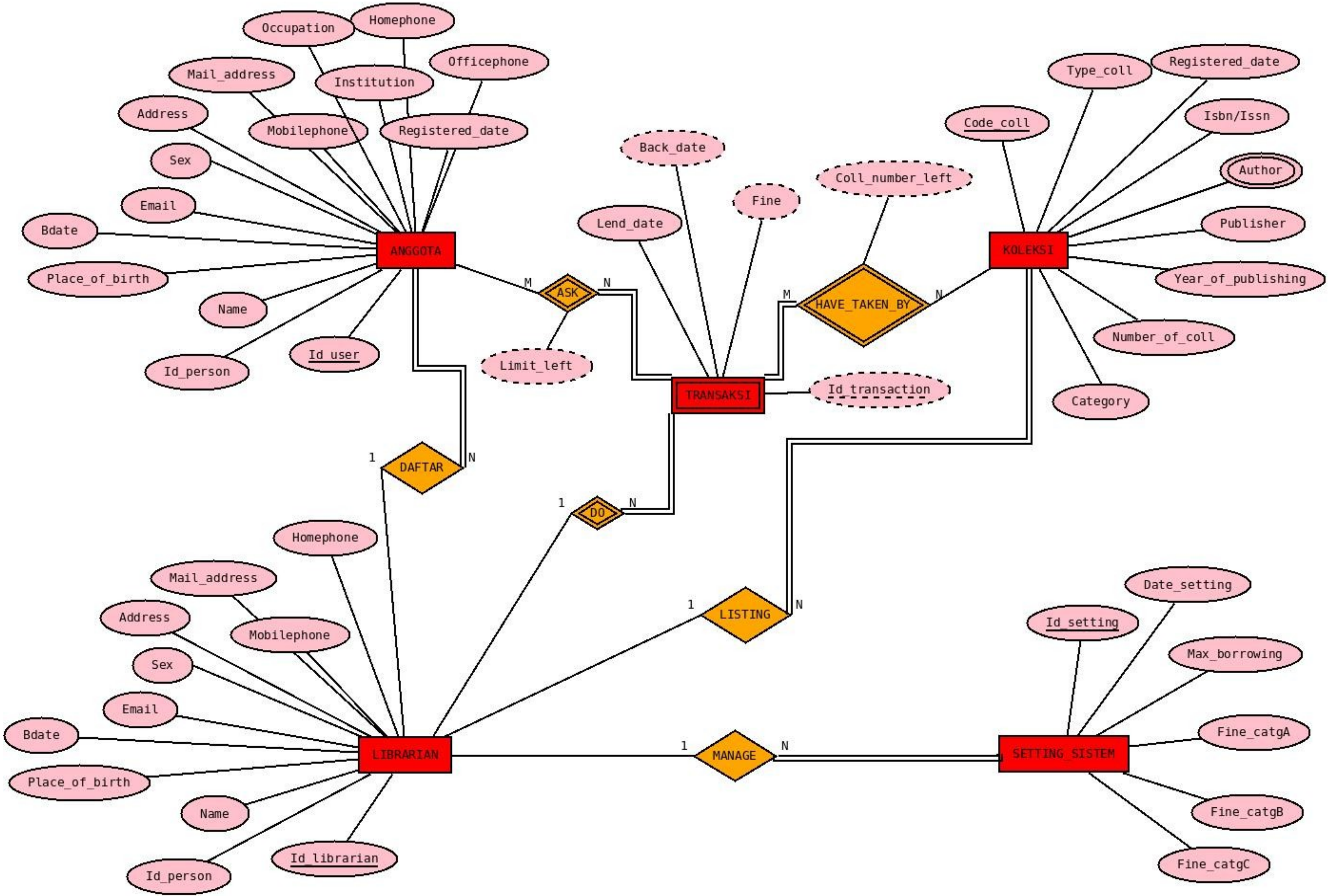
TRUCK

VehicleId	TMake	TModel	Tonnage	TYear

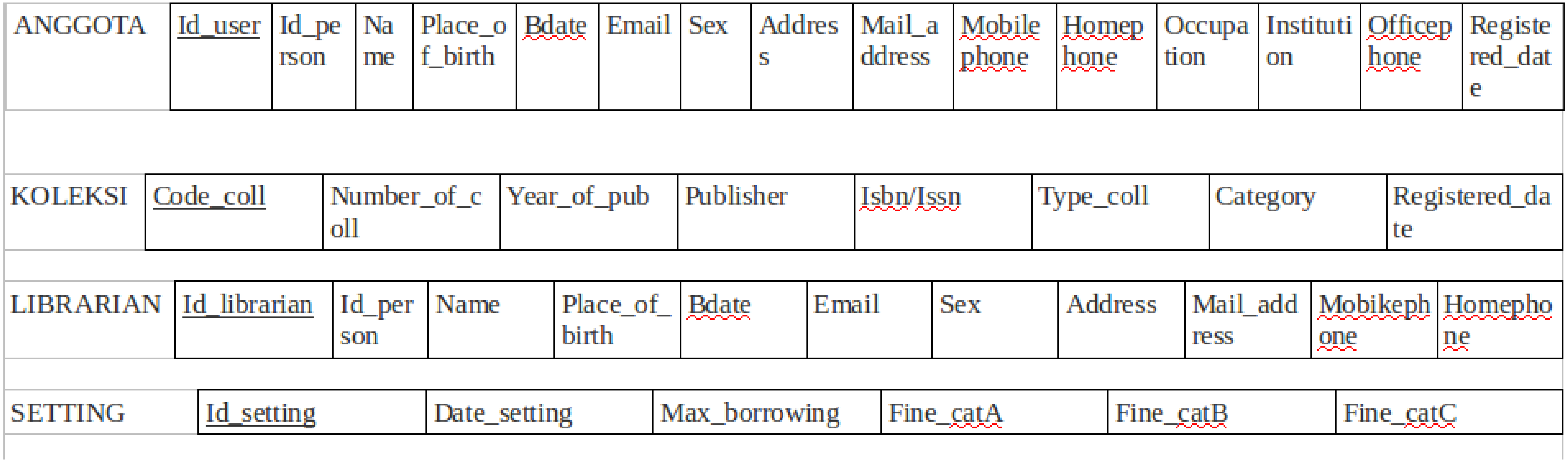
OWNS

OwnerId	VehicleId	PurchaseDate	LienOrRegular

Mapping Exercise 1



Step 1



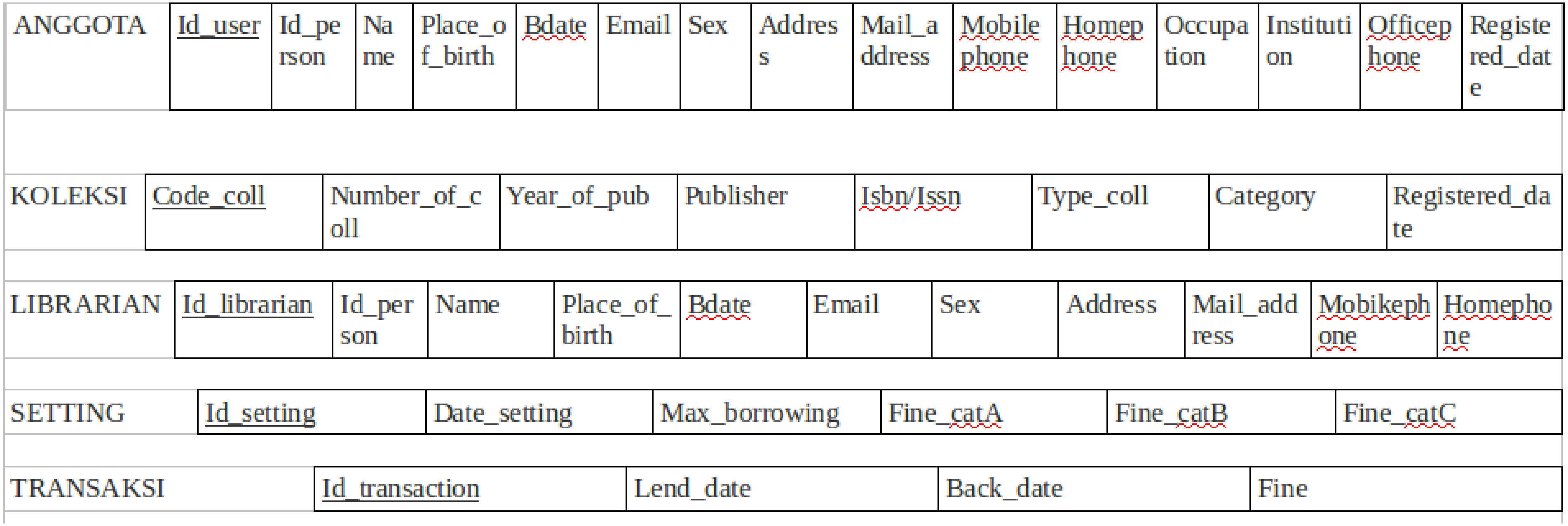
ANGGOTA	Id_user	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobilephone	Homephone	Occupation	Institution	Officephone	Registered_date

KOLEKSI	Code_coll	Number_of_coll	Year_of_pub	Publisher	Isbn/Issn	Type_coll	Category	Registered_date

LIBRARIAN	Id_librarian	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobikephone	Homephone

SETTING	Id_setting	Date_setting	Max_borrowing	Fine_catA	Fine_catB	Fine_catC

Step 2



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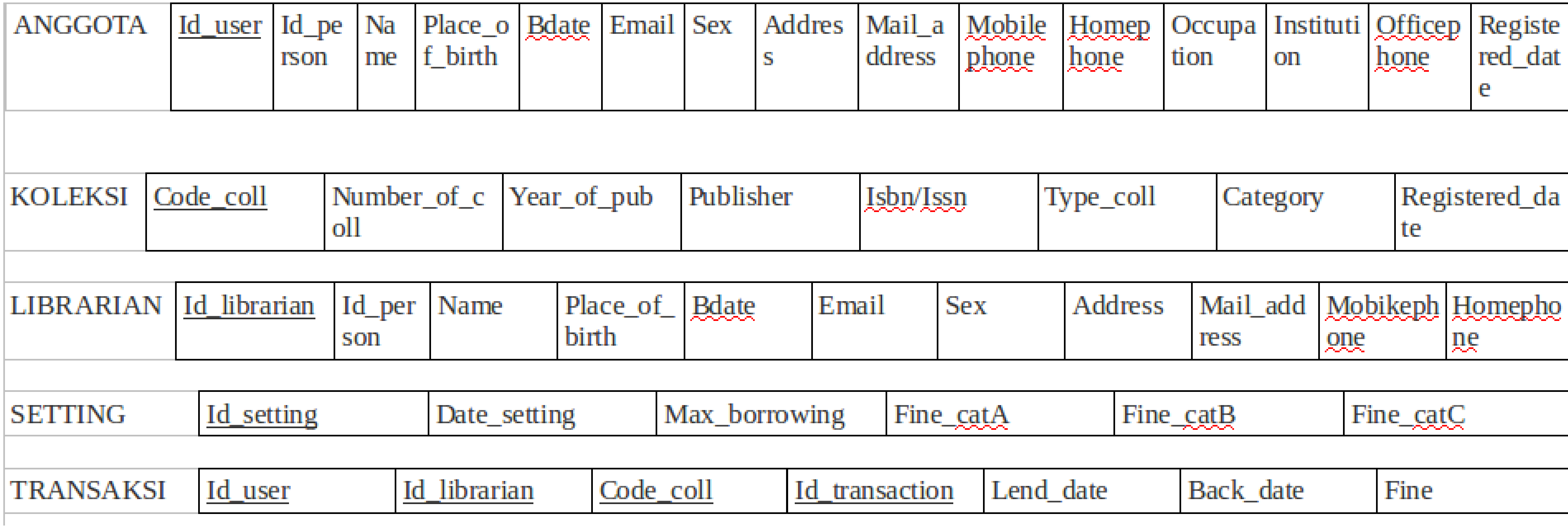
KOLEKSI	Code_coll	Number_of_coll	Year_of_pub	Publisher	Isbn/Issn	Type_coll	Category	Registered_date

LIBRARIAN	Id_librarian	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobikephone	Homephone

SETTING	Id_setting	Date_setting	Max_borrowing	Fine_catA	Fine_catB	Fine_catC

TRANSAKSI	Id_transaction	Lend_date	Back_date	Fine

Step 3, 4



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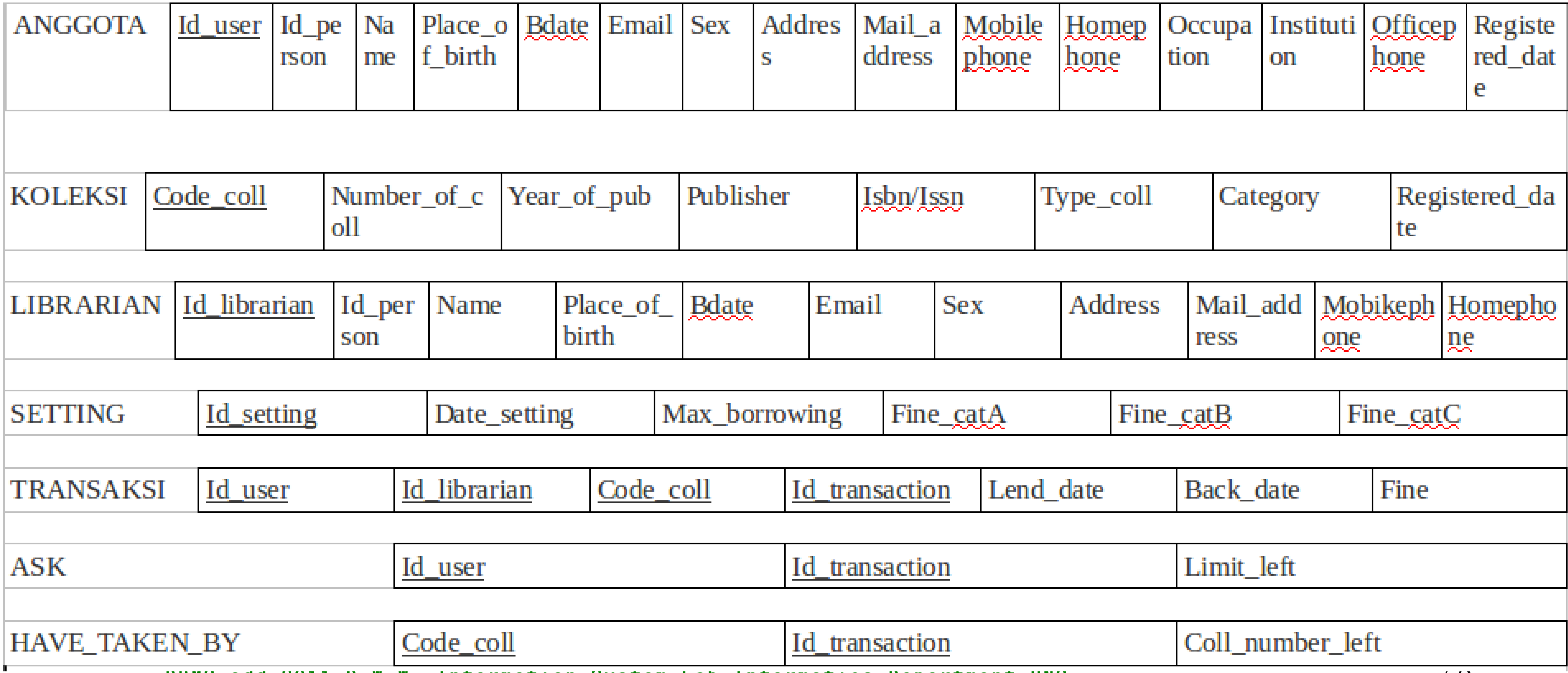
KOLEKSI	Code_coll	Number_of_coll	Year_of_pub	Publisher	Isbn/Issn	Type_coll	Category	Registered_date

LIBRARIAN	Id_librarian	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobikephone	Homephone

SETTING	Id_setting	Date_setting	Max_borrowing	Fine_catA	Fine_catB	Fine_catC

TRANSAKSI	Id_user	Id_librarian	Code_coll	Id_transaction	Lend_date	Back_date	Fine

Step 5



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KOLEKSI	Code_coll	Number_of_coll	Year_of_pub	Publisher	Isbn/Issn	Type_coll	Category	Registered_date

LIBRARIAN	Id_librarian	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobikephone	Homephone

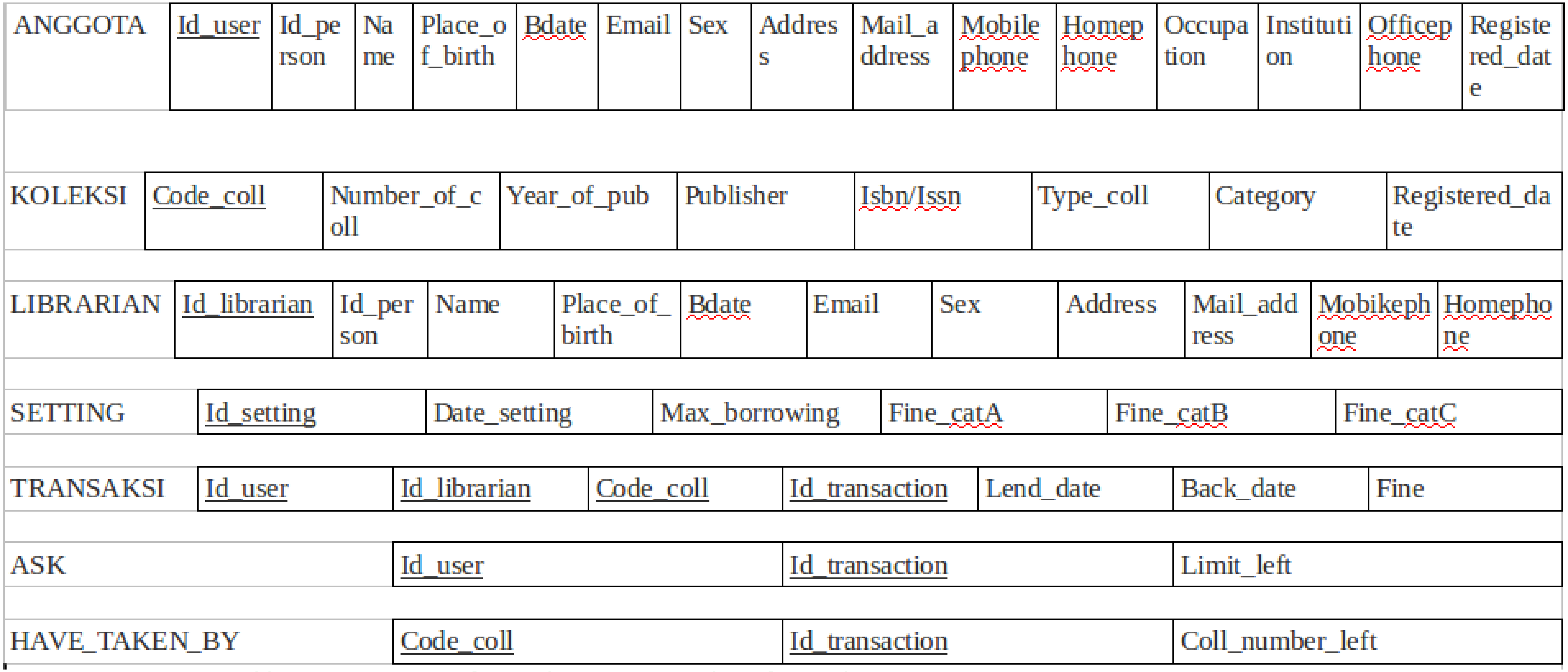
SETTING	Id_setting	Date_setting	Max_borrowing	Fine_catA	Fine_catB	Fine_catC

TRANSAKSI	Id_user	Id_librarian	Code_coll	Id_transaction	Lend_date	Back_date	Fine

ASK	Id_user	Id_transaction	Limit_left

HAVE_TAKEN_BY	Code_coll	Id_transaction	Coll_number_left

Step 6



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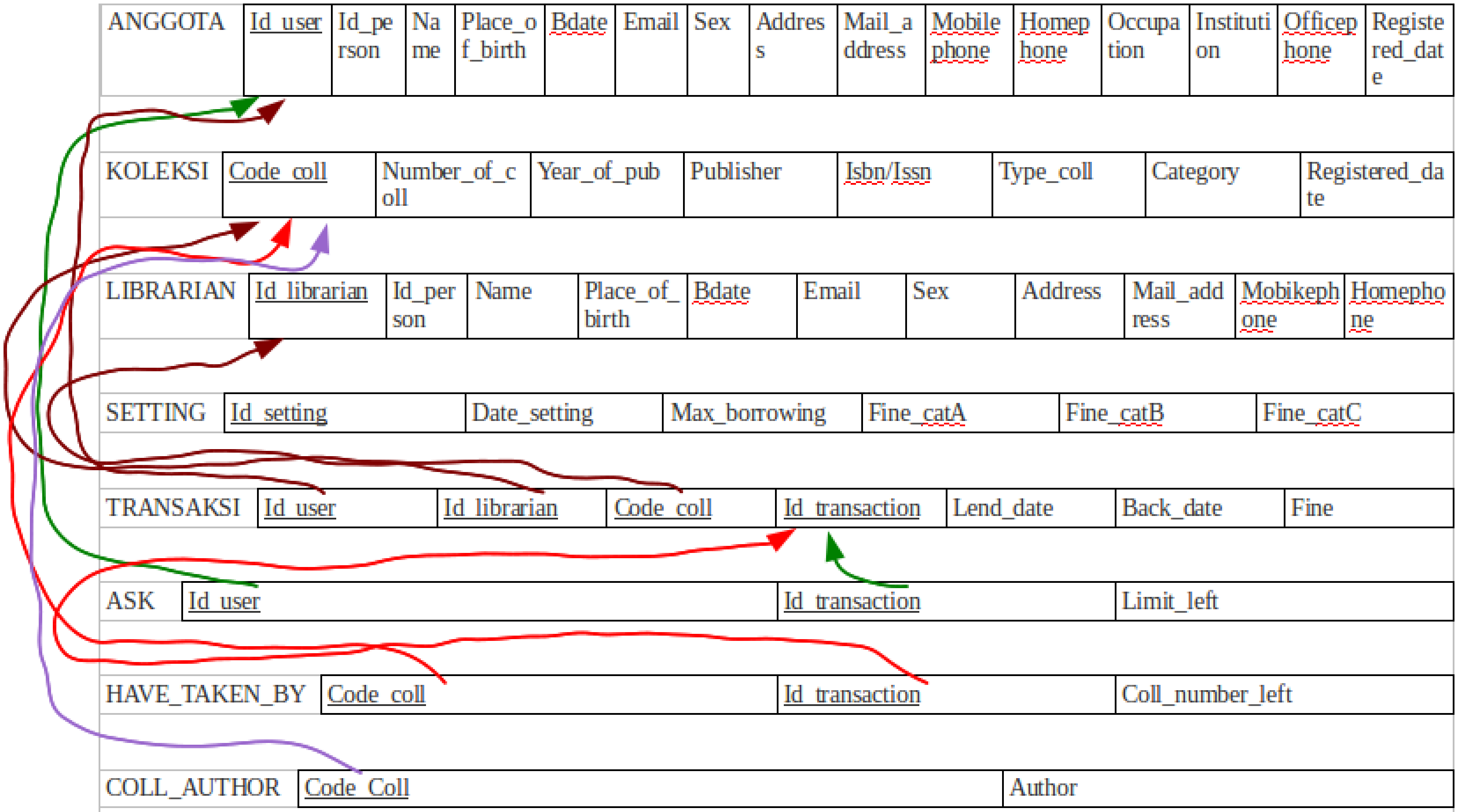
KOLEKSI	Code_coll	Number_of_coll	Year_of_pub	Publisher	Isbn/Issn	Type_coll	Category	Registered_date							
LIBRARIAN	Id_librarian	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobikephone	Homephone				
SETTING	Id_setting	Date_setting	Max_borrowing	Fine_catA	Fine_catB	Fine_catC									
TRANSAKSI	Id_user	Id_librarian	Code_coll	Id_transaction	Lend_date	Back_date	Fine								
ASK	Id_user	Id_transaction	Limit_left												
HAVE_TAKEN_BY	Code_coll	Id_transaction	Coll_number_left												

. . .



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KOLEKSI	Code_coll	Number_of_coll	Year_of_pub	Publisher	Isbn/Issn	Type_coll	Category	Registered_date

LIBRARIAN	Id_librarian	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobikephone	Homephone

SETTING	Id_setting	Date_setting	Max_borrowing	Fine_catA	Fine_catB	Fine_catC

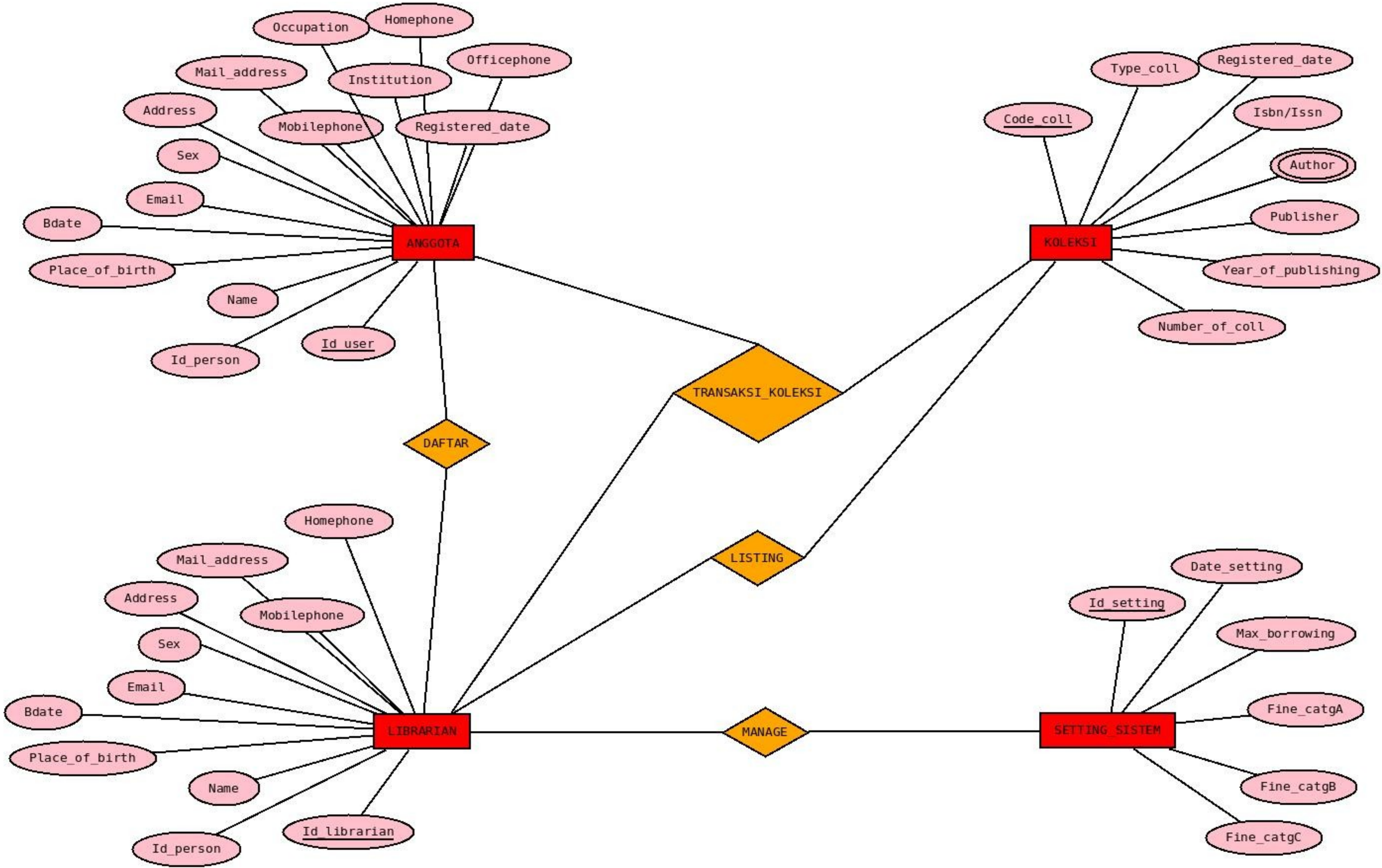
TRANSAKSI	Id_user	Id_librarian	Code_coll	Id_transaction	Lend_date	Back_date	Fine

ASK	Id_user	Id_transaction	Limit_left

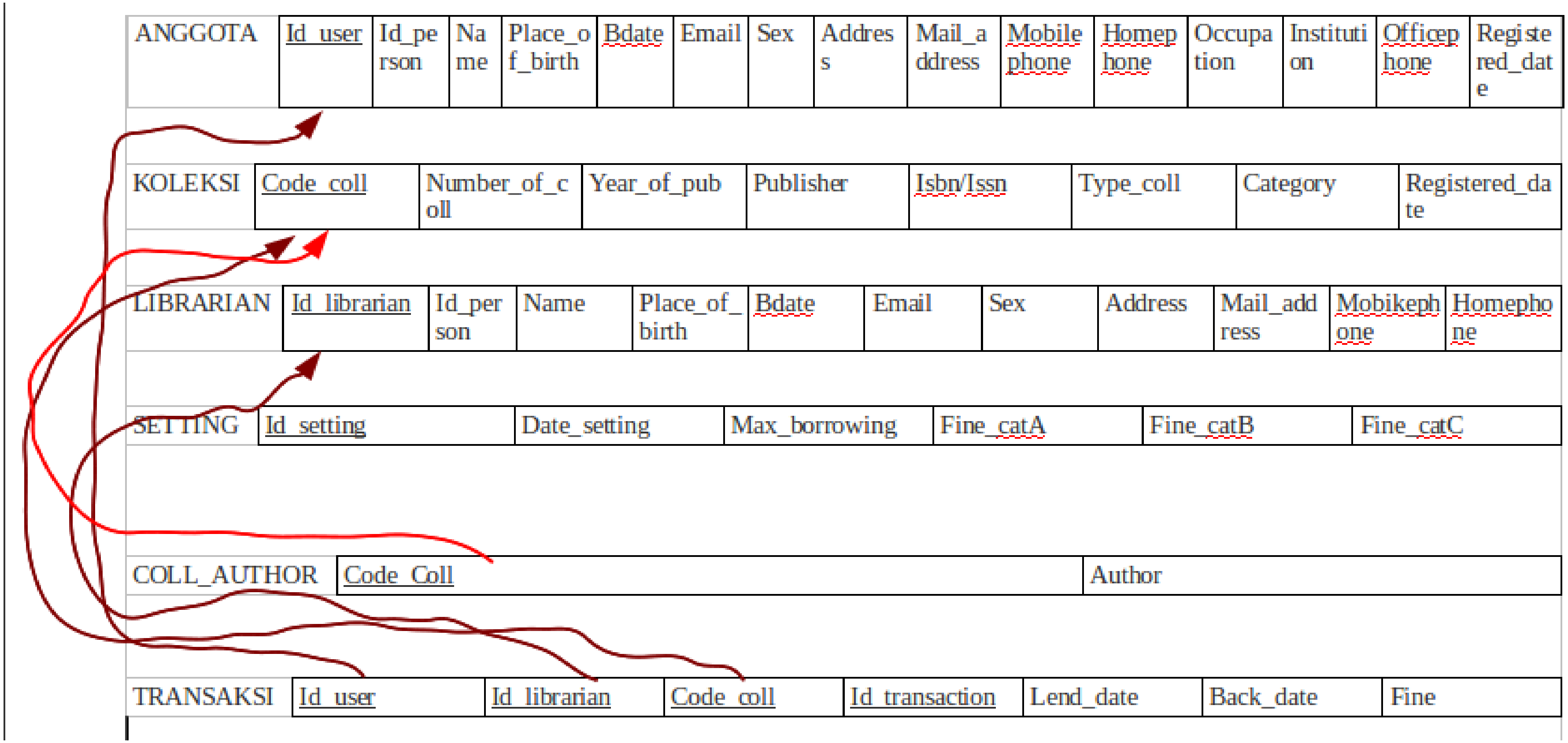
HAVE_TAKEN_BY	Code_coll	Id_transaction	Coll_number_left

COLL_AUTHOR	Code_Coll	Author

What if this ER



We will have step 7 … :)



ANGGOTA	Id_user	Id_person	Name	Place_of_birth	Bdate	Email	Sex	Address	Mail_address	Mobilephone	Homephone	Occupation	Institution	Officephone	Registered_date

KOLEKSI	Code_coll	Number_of_coll	Year_of_pub	Publisher	Isbn/Issn	Type_coll	Category	Registered_date

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SETTING	Id_setting	Date_setting	Max_borrowing	Fine_catA	Fine_catB	Fine_catC

COLL_AUTHOR	Code_Coll	Author

TRANSAKSI	Id_user	Id_librarian	Code_coll	Id_transaction	Lend_date	Back_date	Fine